

# Adversarial Search

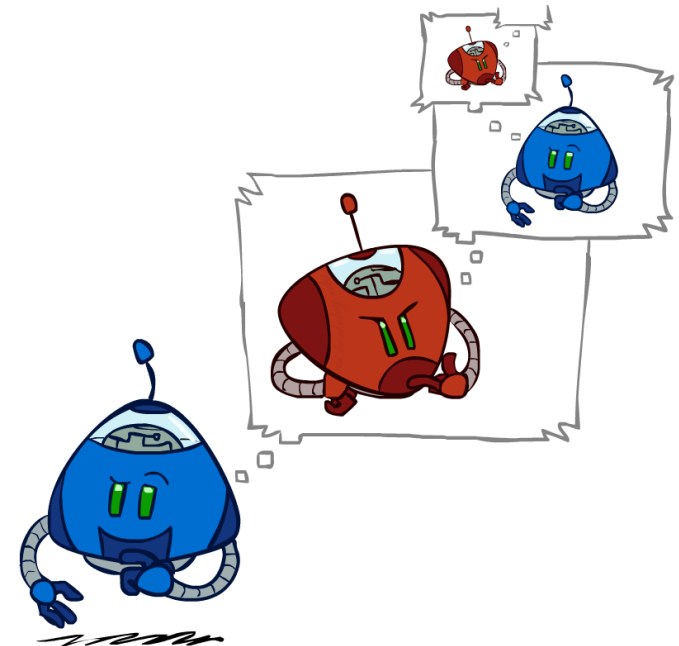
## Lecture 7

How can we use **search to plan ahead** when **other agents are planning against us?**



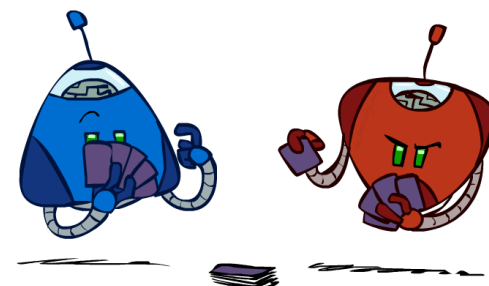
# Agenda

- Games: context, history
- Searching via Minimax
- Scaling
  - $\alpha$ – $\beta$  pruning
  - Depth-limiting
  - Evaluation functions
- Handling uncertainty with Expectiminimax



# Characterizing Games

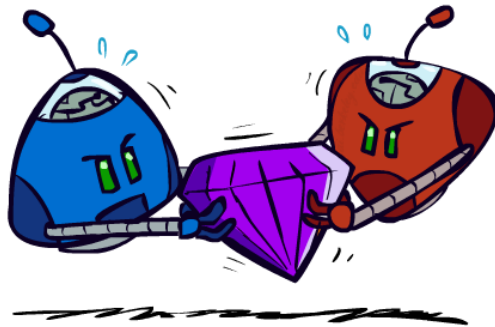
- There are many kinds of games, and several ways to classify them
  - Deterministic vs. stochastic
  - [Im]perfect information
  - One, two, multi-player
  - Utility (how agents value outcomes)
    - **Zero-sum**
- Algorithmic goal: calculate a **strategy** (or **policy**) that decides a move in each state



# Utility

## Zero/Constant-Sum

- Opposite utilities
- Adversarial, pure competition



## General Games

- Independent utilities
- Cooperation, indifference, competition, and more are all possible



# Examples: Perception vs. Chance

	Deterministic	Stochastic
Perfect	Chess, Checkers, Go, Othello	Backgammon, Monopoly
Imperfect	Battleship	Bridge, Poker, Scrabble



# Checkers

- 1950: First computer player
- 1994: First computer champion (Chinook) ended 40-year-reign of human champion Marion Tinsley using complete 8-piece endgame
- 1995: defended against Don Lafferty
- 2007: **solved!**



# Chess

- 1997: Deep Blue defeats human champion Gary Kasparov in a six-game match
- Deep Blue examined 200M positions per second, used very sophisticated evaluation and undisclosed methods for extending some lines of search up to 40 ply
- Current programs are *even better*, if less historic

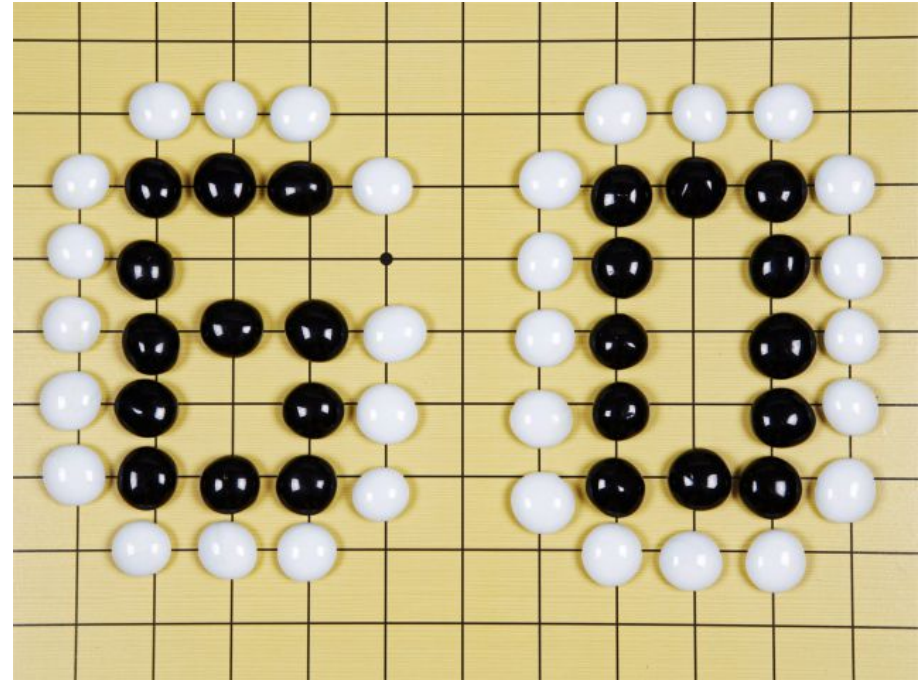


DeepBlue



# Go

- Until recently, AI was not competitive at champion level
  - 2015: beat Fan Hui, European champion (2-dan; 5-0)
  - 2016: beat Lee Sedol, one of the best players in the world (9-dan; 4-1)
  - 2017: beat Ke Jie, #1 in the world (9-dan; 2-0)
- MCTS + ANNs for **policy** (what to do) and **evaluation** (how good is a board state)





# Poker

- Libratus beat four top-class human poker players in January, 2017
  - 120,000 hands played
- Novel methods for endgame solving in imperfect games
- 15 million core hours of computation (+4 during competition)

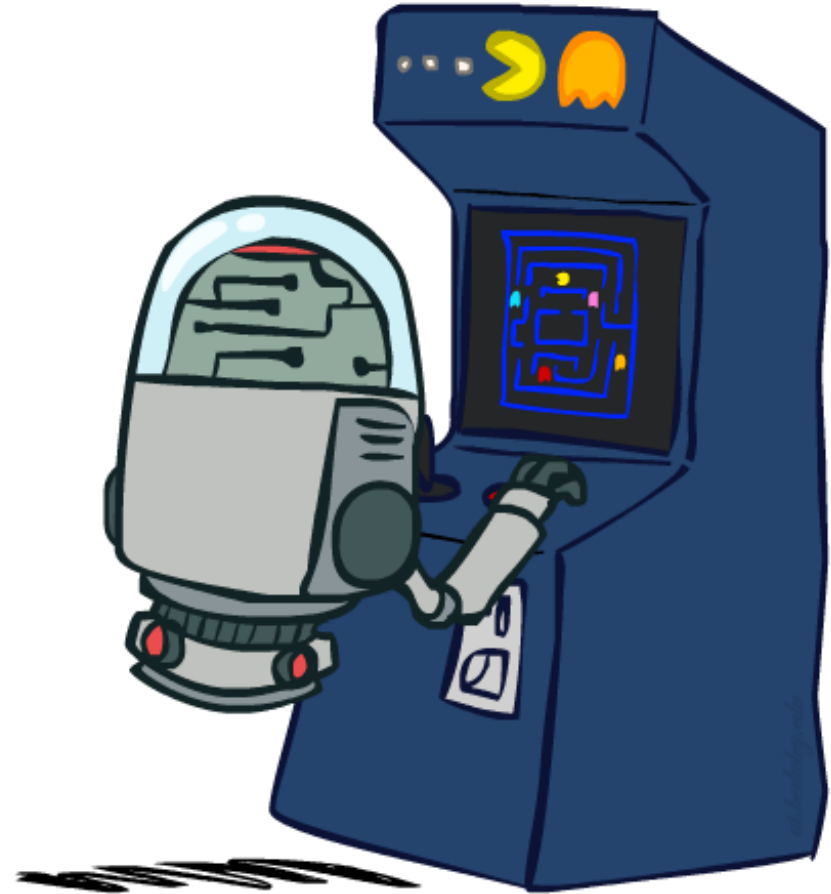


Libratus



# More Progress

- Othello: 1997, defeated world champion
- Bridge: 1998, competitive with human champions
- Scrabble: 2006, defeated world champion



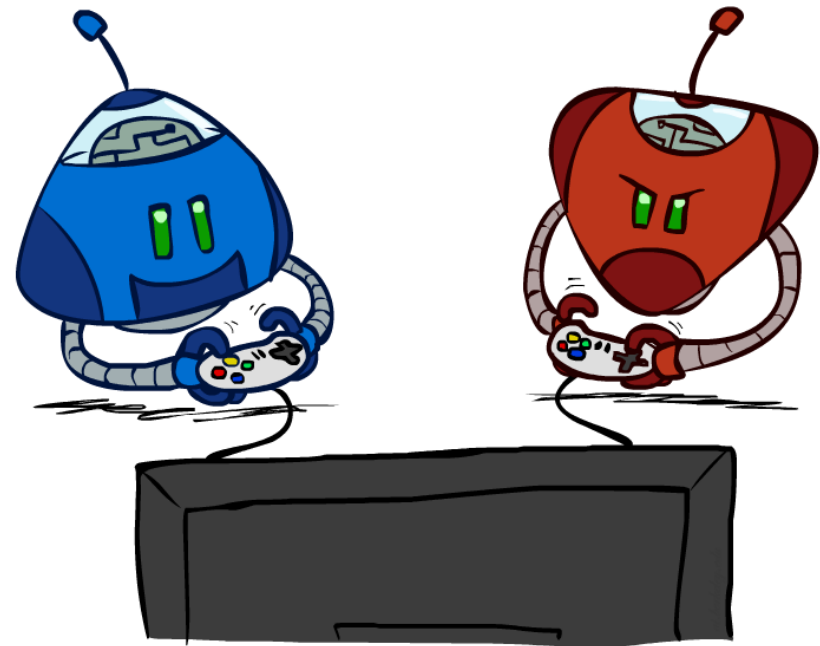
# Game Formalism

- States:  $S$  (start at  $S_0$ )
- Players:  $P \{1, \dots, N\}$  (typically take turns)
- Actions:  $Action(s)$ , returns legal options
- Transition function:  $S \times A \rightarrow S$
- Terminal test:  $Terminal(s)$ , returns T/F
- Utility:  $S \times P \rightarrow \mathbb{R}$
  
- Solution for a player is a **policy**:  $S \rightarrow A$

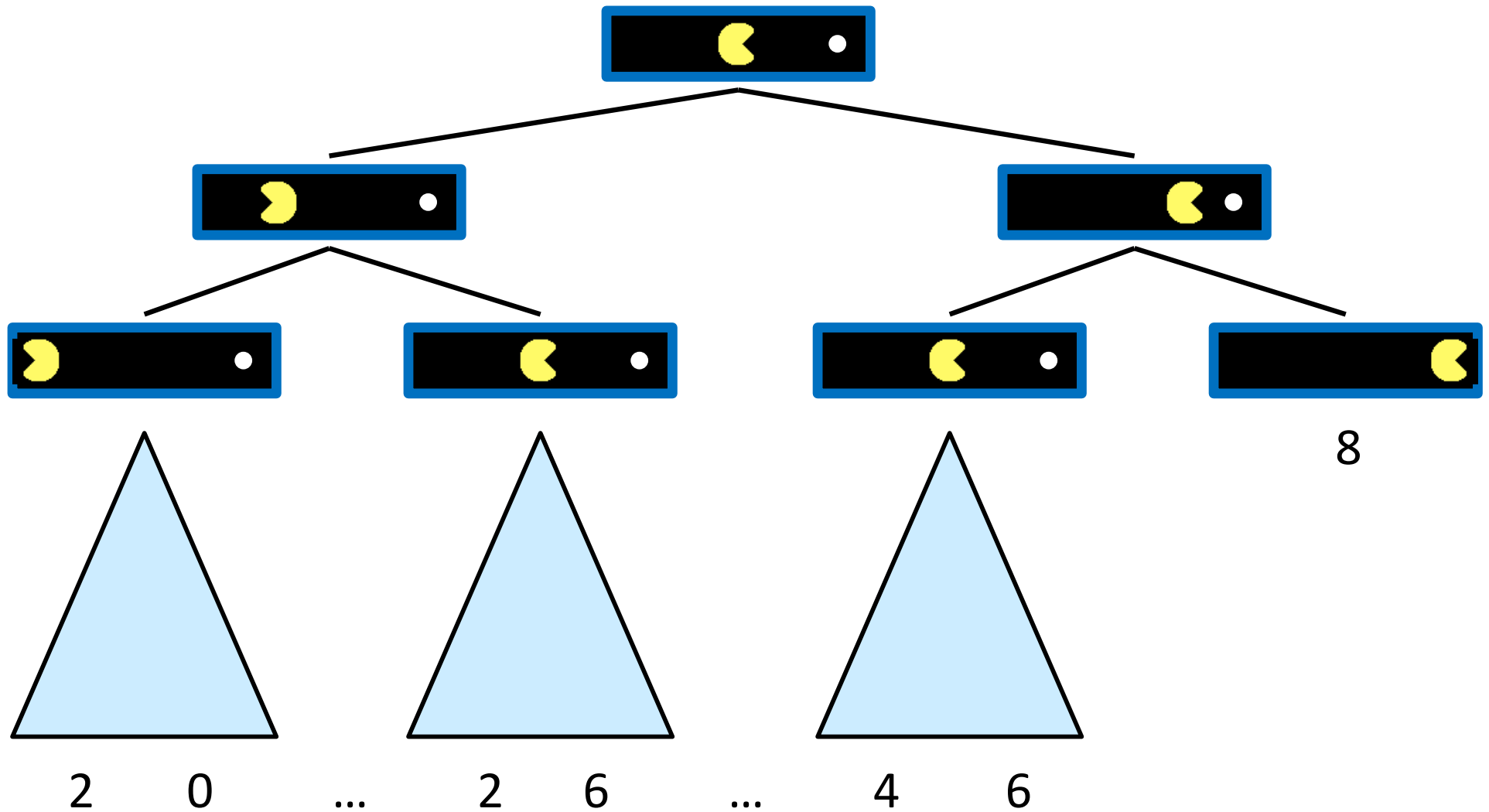


# Game Plan :)

- Start with deterministic, two-player adversarial games
- Issues to come
  - Multiple players
  - Resource limits
  - Stochasticity



# Single-Agent Game Tree

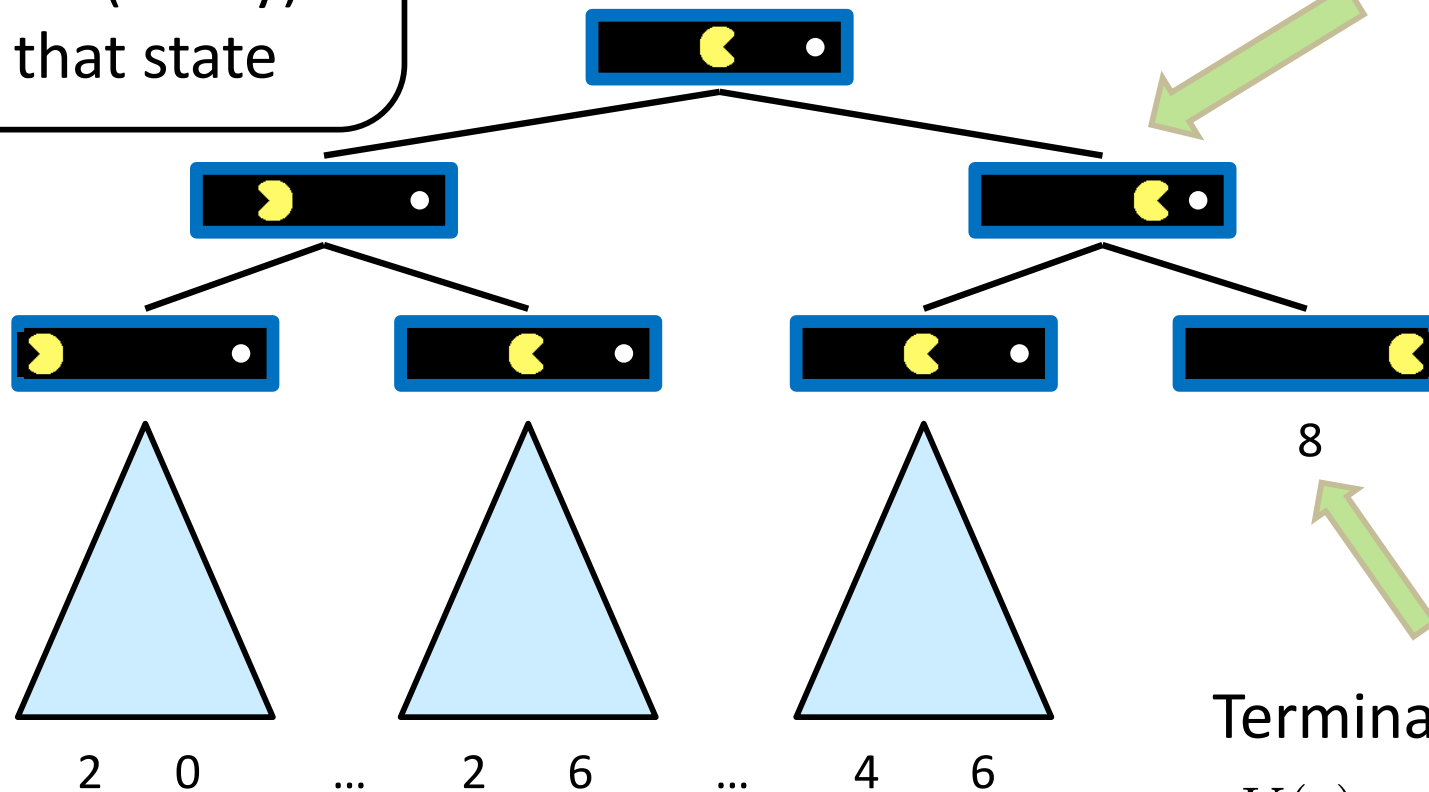


# Value of a State

Value of a state:  
The best achievable  
outcome (utility)  
from that state

Non-Terminal States:

$$V(s) = \max_{s' \in \text{children}(s)} V(s')$$

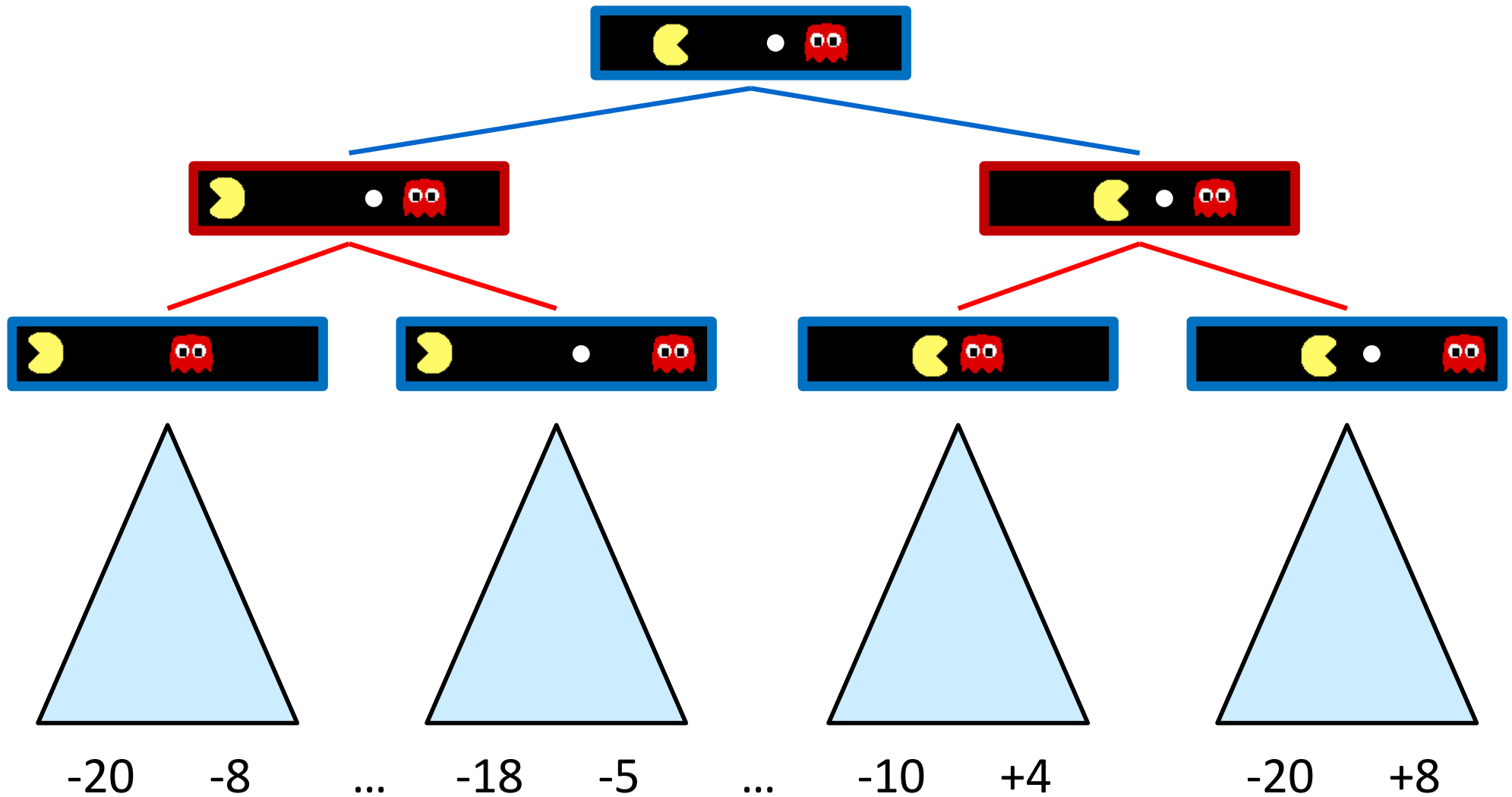


Terminal States:

$$V(s) = \text{known}$$



# Adversarial Game Trees



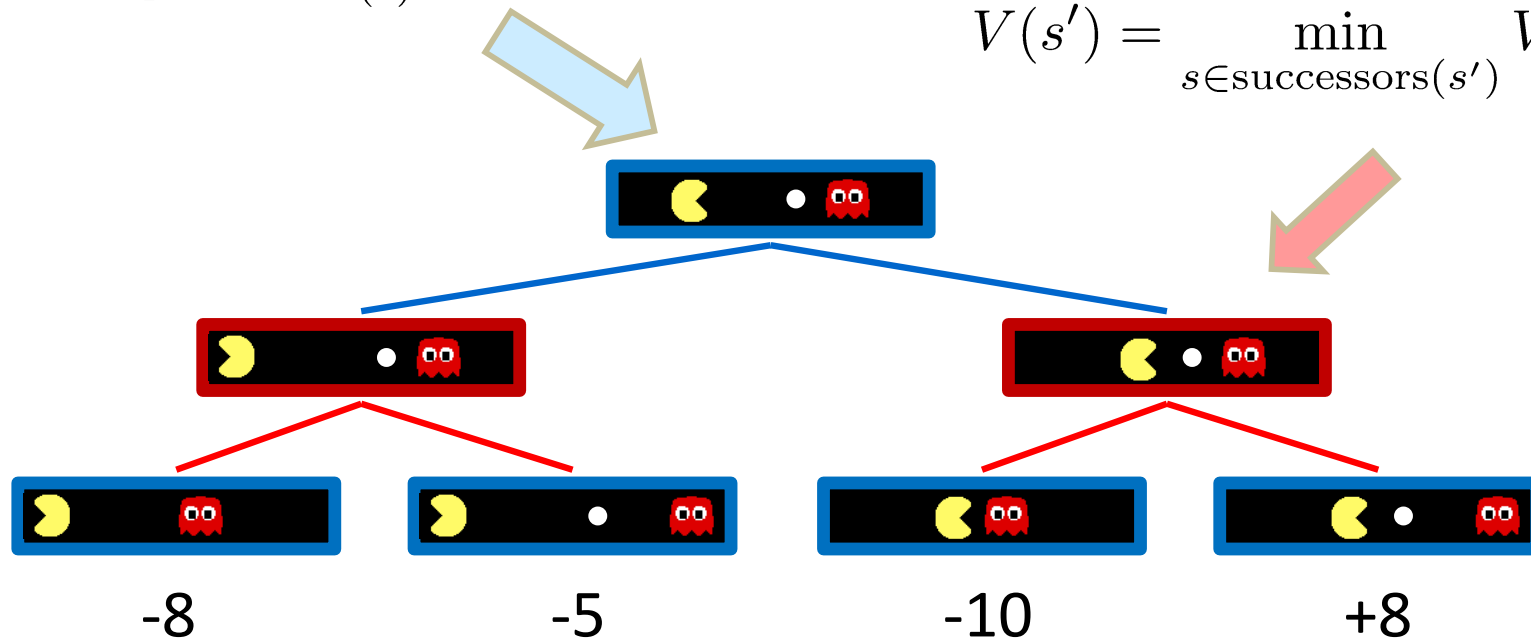
# Minimax Values

States Under Agent's Control:

$$V(s) = \max_{s' \in \text{successors}(s)} V(s')$$

States Under Opponent's Control:

$$V(s') = \min_{s \in \text{successors}(s')} V(s)$$



Terminal States:

$$V(s) = \text{known}$$





# Tic-Tac-Toe Game Tree



MAX (X)



MIN (O)



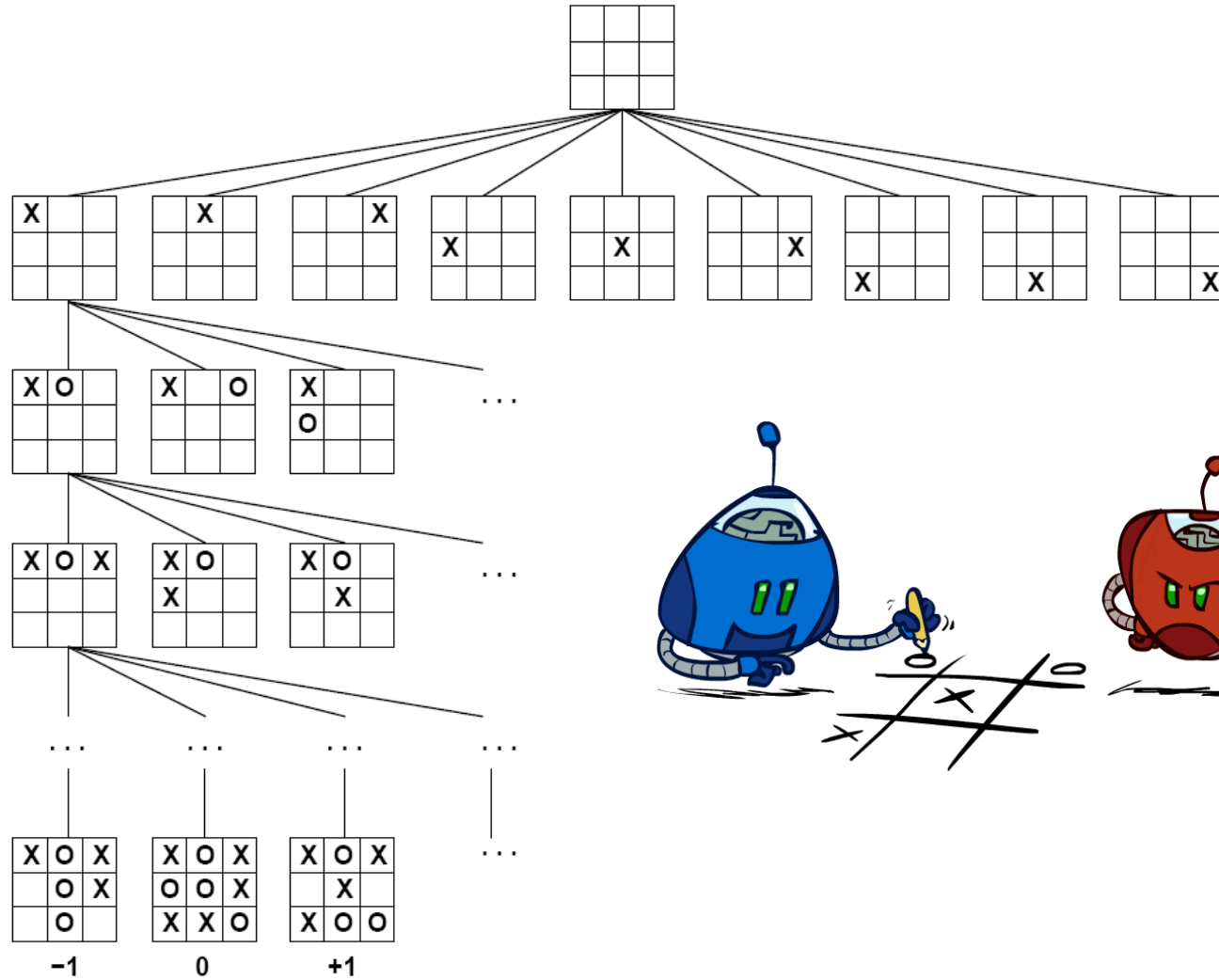
MAX (X)



MIN (O)

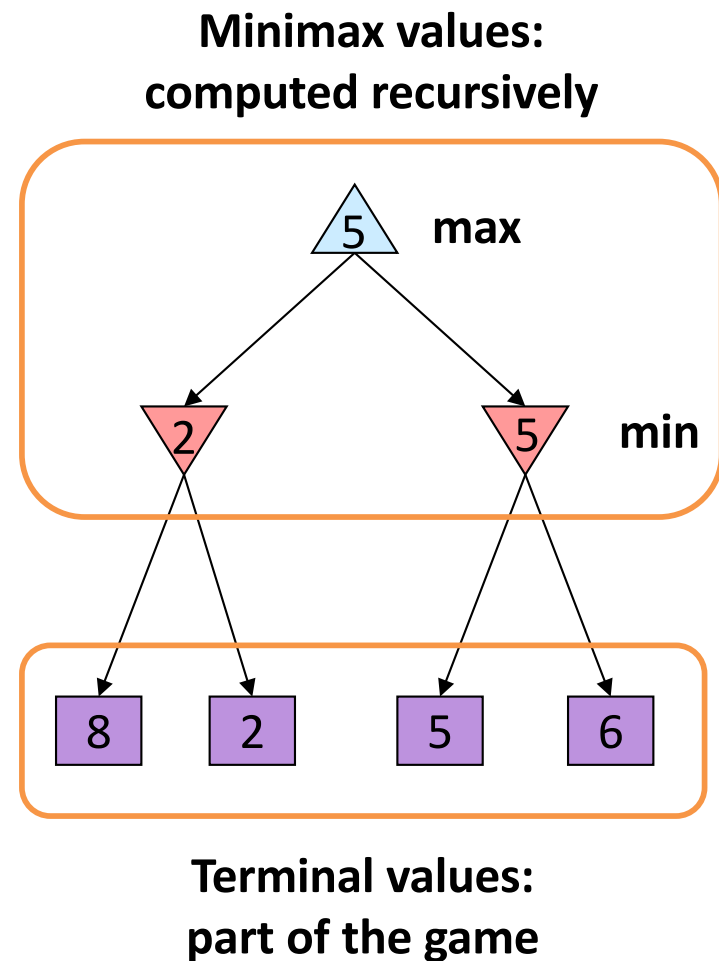
TERMINAL

Utility



# Adversarial Search via Minimax

- Deterministic, zero-sum
  - Tic-tac-toe, chess
  - One player maximizes
  - The other minimizes
- Minimax search
  - A search tree
  - Players alternate turns
  - Compute each node's *minimax value*: **the best achievable utility against a rational (optimal) adversary**



# Minimax Implementation

```
def value(state):  
    if the state is a terminal state: return the state's utility  
    if the next agent is MAX: return max-value(state)  
    if the next agent is MIN: return min-value(state)
```



```
def max-value(state):  
    initialize v =  $-\infty$   
    for each successor of state:  
        v = max(v, value(successor))  
    return v
```



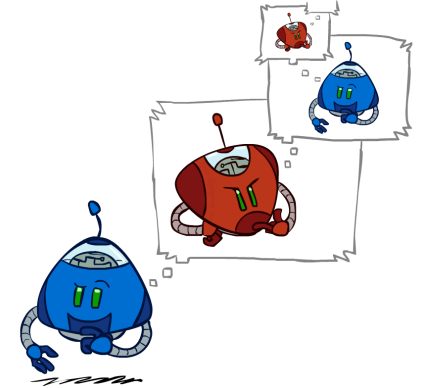
```
def min-value(state):  
    initialize v =  $+\infty$   
    for each successor of state:  
        v = min(v, value(successor))  
    return v
```



# Minimax Evaluation

## Time

- $O(b^m)$ 
  - For chess:  
 $b \approx 35, m \approx 100$

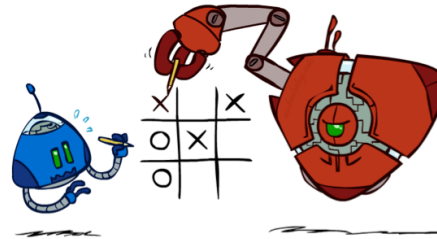


## Space

- $O(bm)$

## Complete

- Only if finite

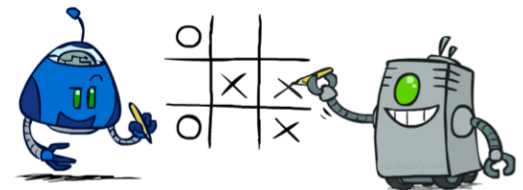


Minimax-Min

## Optimal

- Yes, against **optimal** opponent

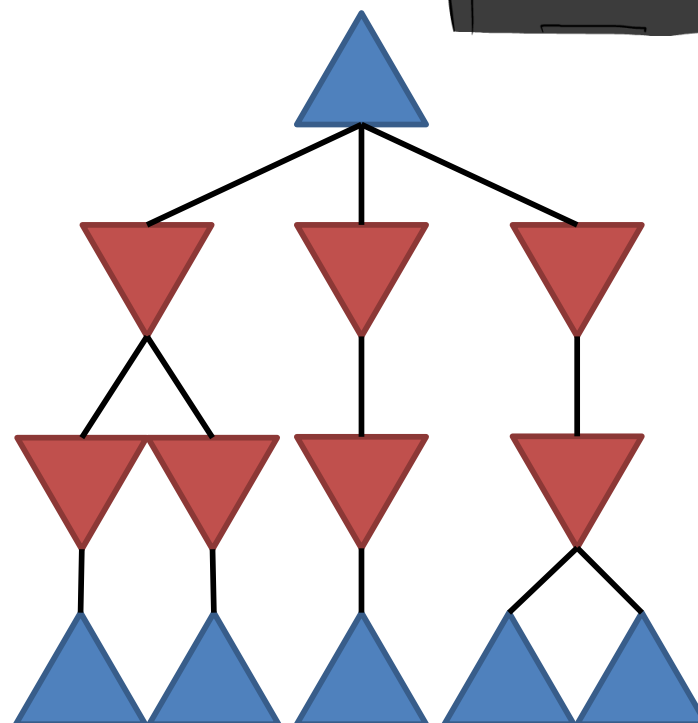
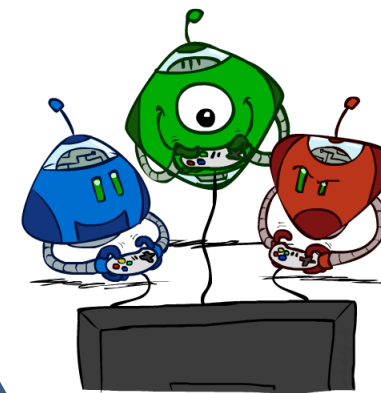
Minimax-Avg



# Multiple Players

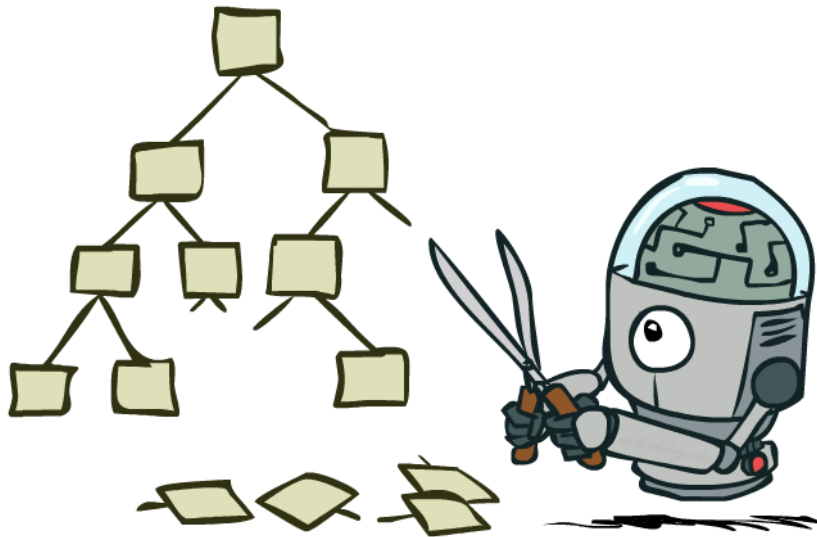
Add a **ply** per player

- Independent utility: use a vector of values, each player MAX own utility
- Zero-sum: each team sequentially MIN/MAX
  - In Pacman, have multiple MIN layers for each ghost per 1 Pacman move

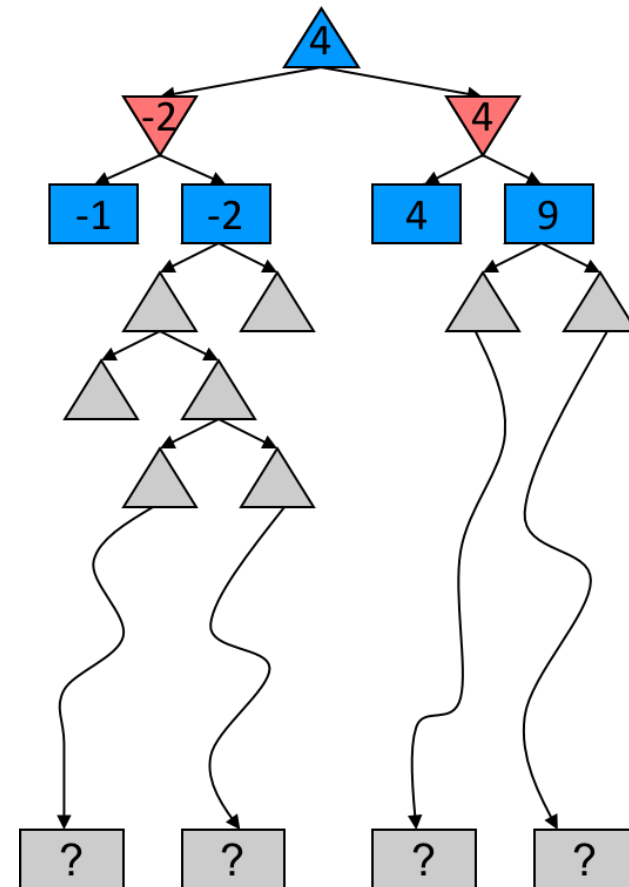


# Scaling to Larger Games

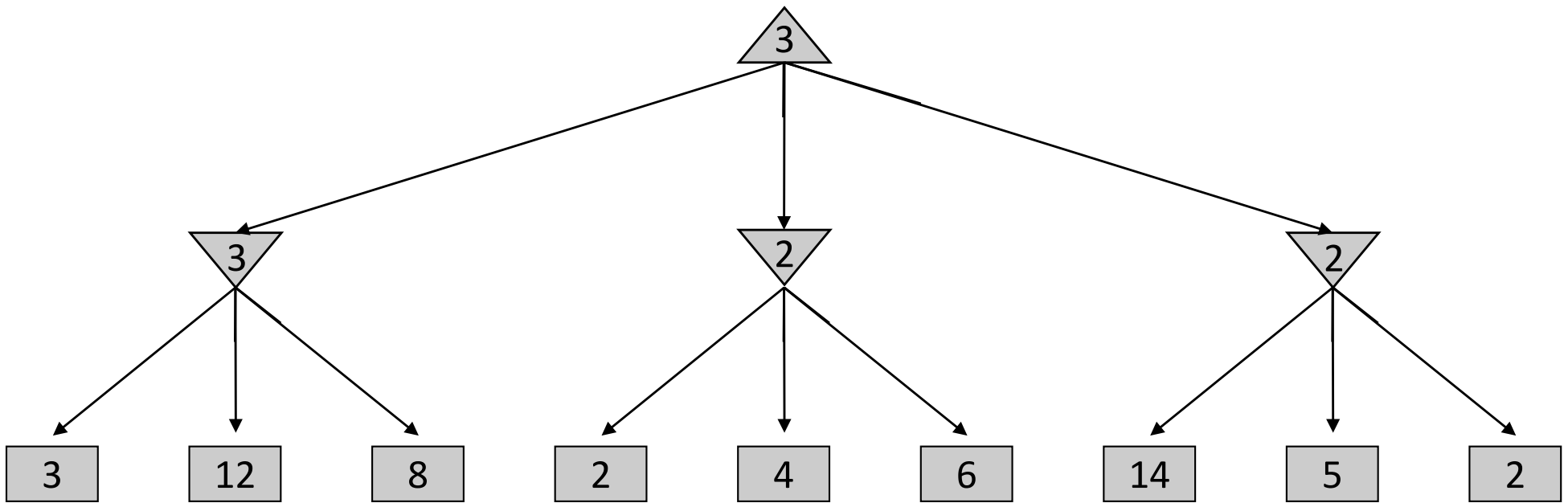
## Tree Pruning



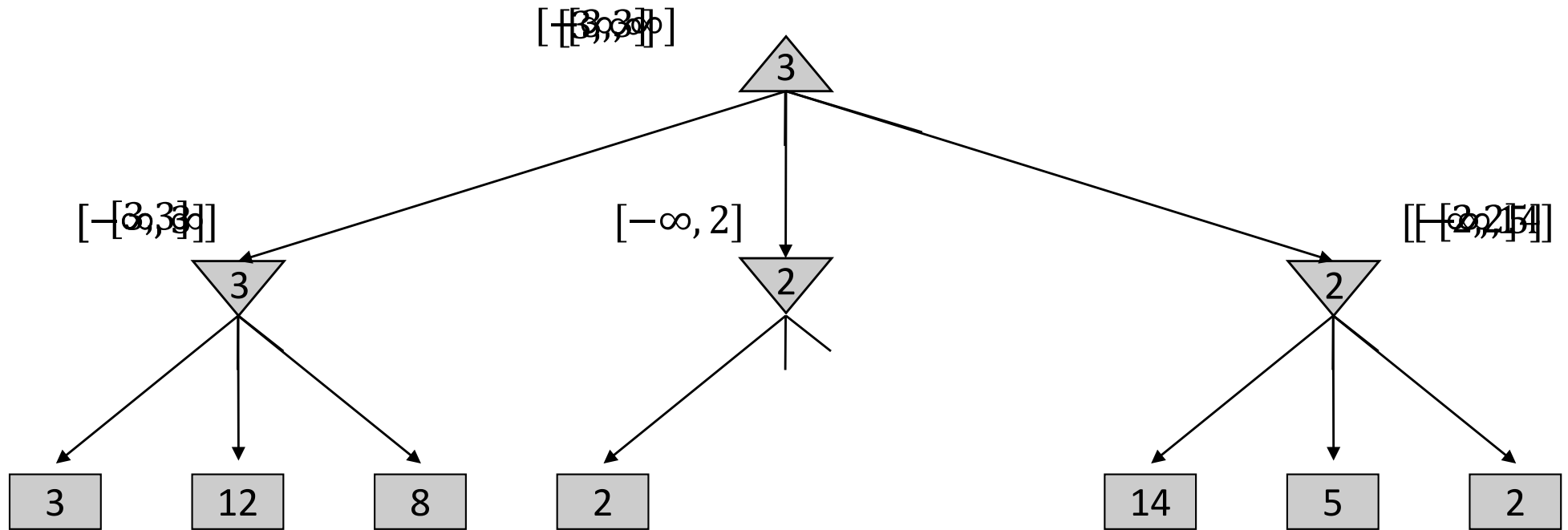
## Depth-Limiting + Evaluation



# Minimax Example

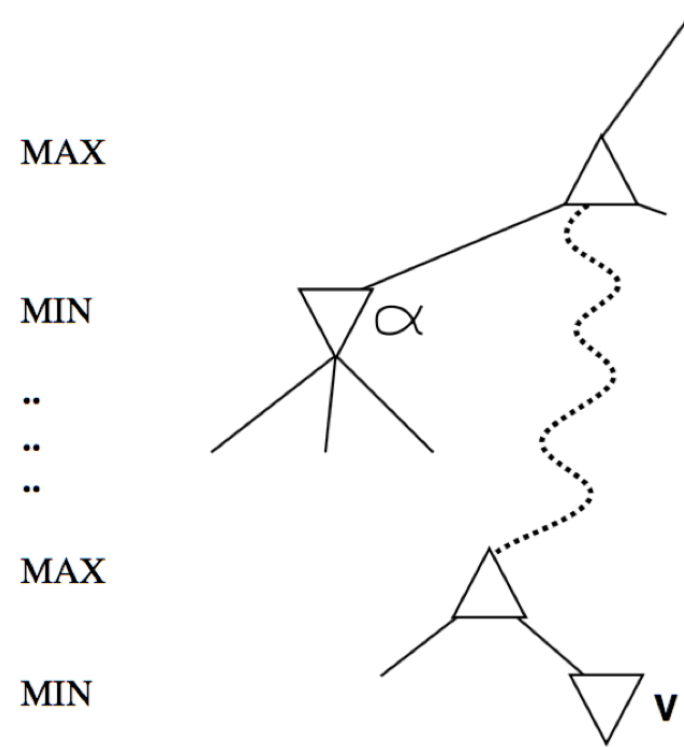


# Minimax Pruning





# General Case



- $\alpha$  is the best value (to *MAX*) found so far off the current path
- If  $V$  is worse than  $\alpha$ , *MAX* will avoid it – prune that branch
- Define  $\beta$  similarly for *MIN*



# Alpha-Beta Pruning

```
def min-value(state,  $\alpha$ ,  $\beta$ ):  
    initialize  $v = +\infty$   
    for each successor of state:  
         $v = \min(v, \text{value}(\text{successor}, \alpha, \beta))$   
        if  $v \leq \alpha$  return  $v$   
         $\beta = \min(\beta, v)$   
    return  $v$ 
```

$\alpha$ : MAX's best option on path  
 $\beta$ : MIN's best option on path

```
def max-value(state,  $\alpha$ ,  $\beta$ ):  
    initialize  $v = -\infty$   
    for each successor of state:  
         $v = \max(v, \text{value}(\text{successor}, \alpha, \beta))$   
        if  $v \geq \beta$  return  $v$   
         $\alpha = \max(\alpha, v)$   
    return  $v$ 
```

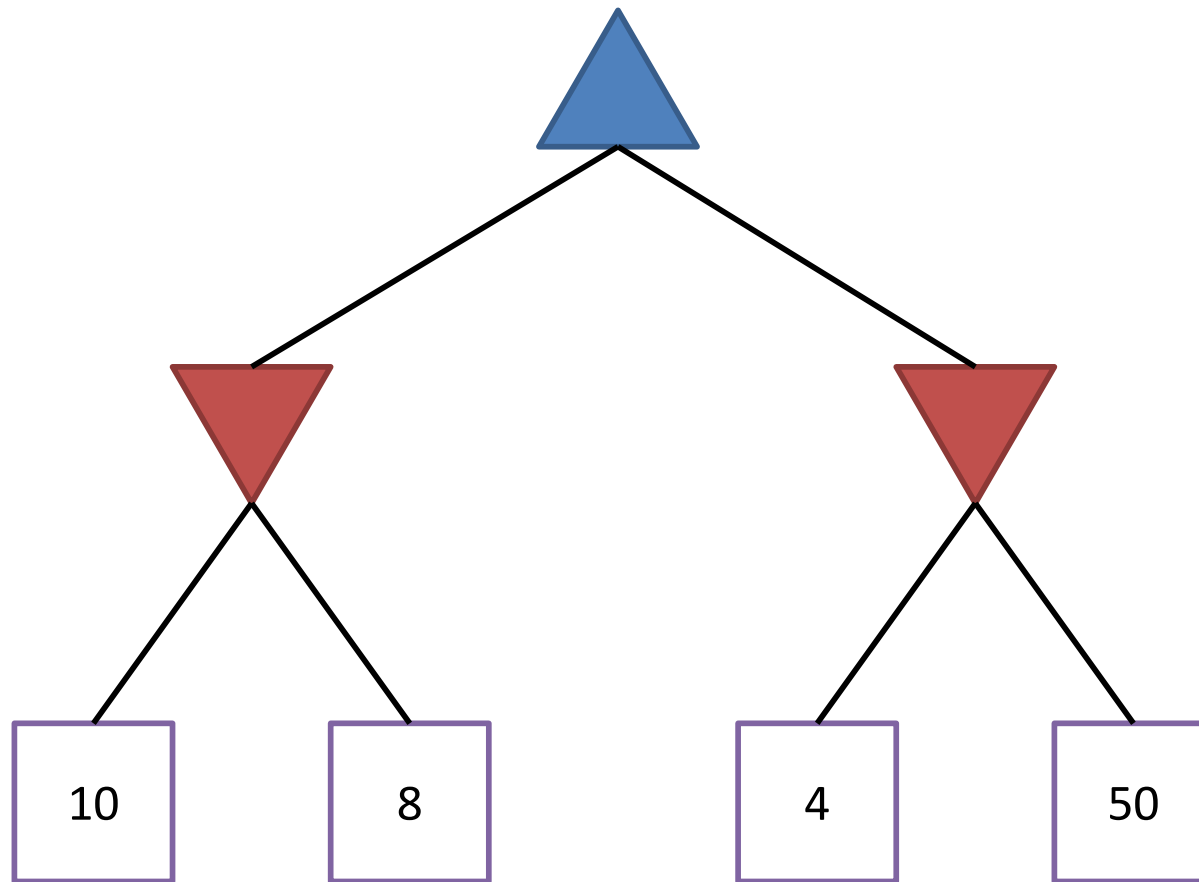


# Alpha-Beta Properties

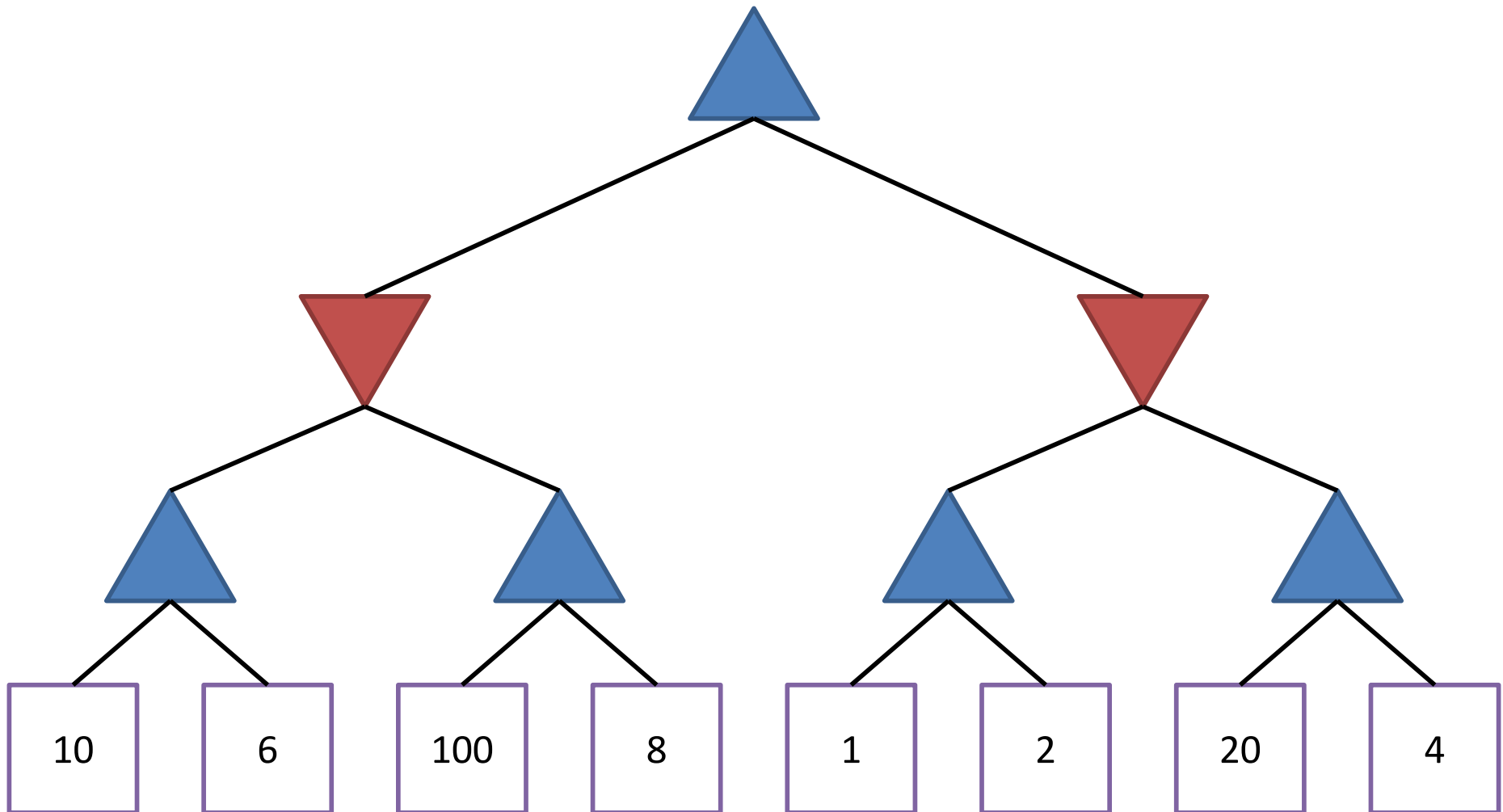
- Has no effect on minimax value computed for the root!
- Good child ordering improves effectiveness of pruning
- With “perfect ordering”:
  - Time complexity drops to  $\mathcal{O}(b^{m/2})$
  - Doubles solvable depth!
  - Full search of, e.g. chess, is still hopeless...
- This is a simple example of **metareasoning** (computing about what to compute)



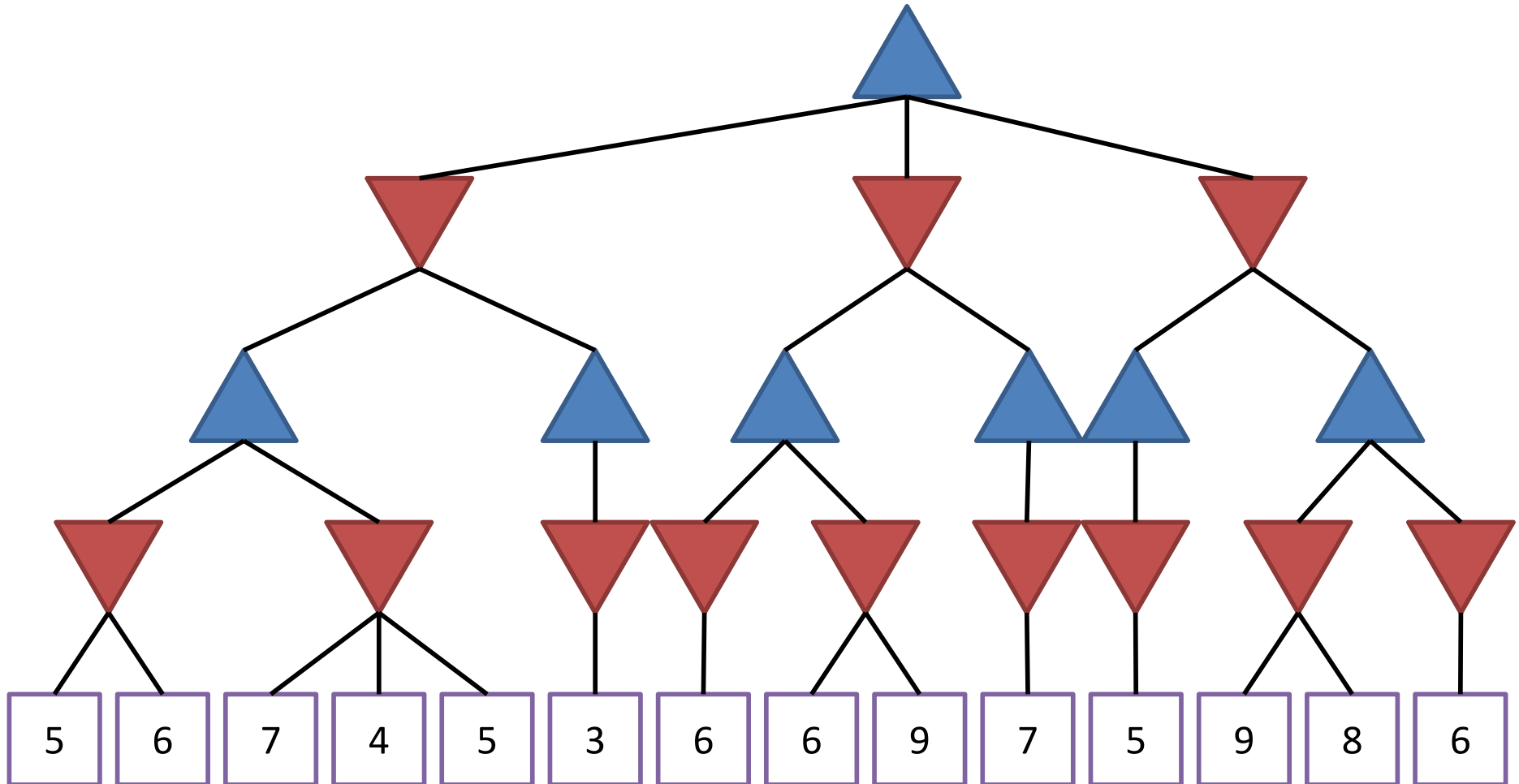
# Checkpoint #1



# Checkpoint #2



# Checkpoint #3



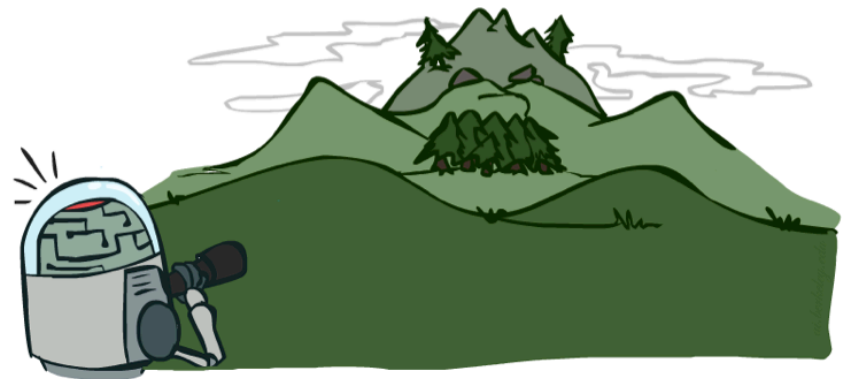
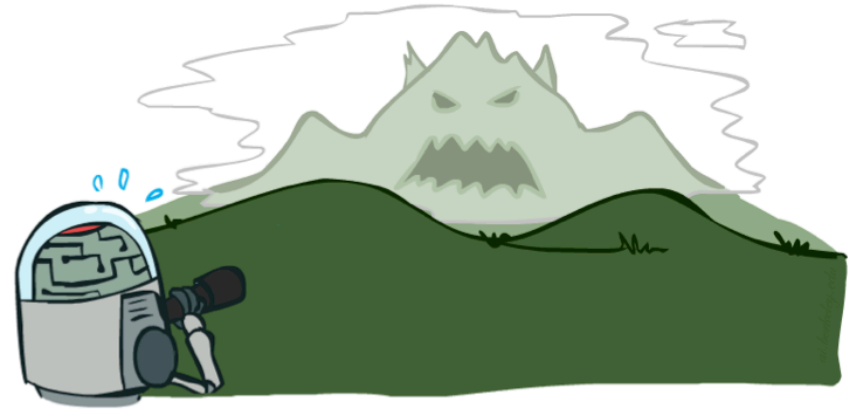
# Resource Limits



- Problem: in realistic games, cannot search to leaves!
- Solution: depth-limited search
  1. Search only to a limited depth in the tree
  2. Replace terminal utilities with an **evaluation function** for non-terminal positions
- Guarantee of optimal play is gone
- More plies makes a BIG difference
- Use iterative deepening for an anytime algorithm



# Search Depth Matters

- Evaluation functions are always imperfect
- The deeper in the tree the evaluation function is buried, the less the quality of the evaluation function matters
- An important example of the tradeoff between complexity of features and complexity of computation

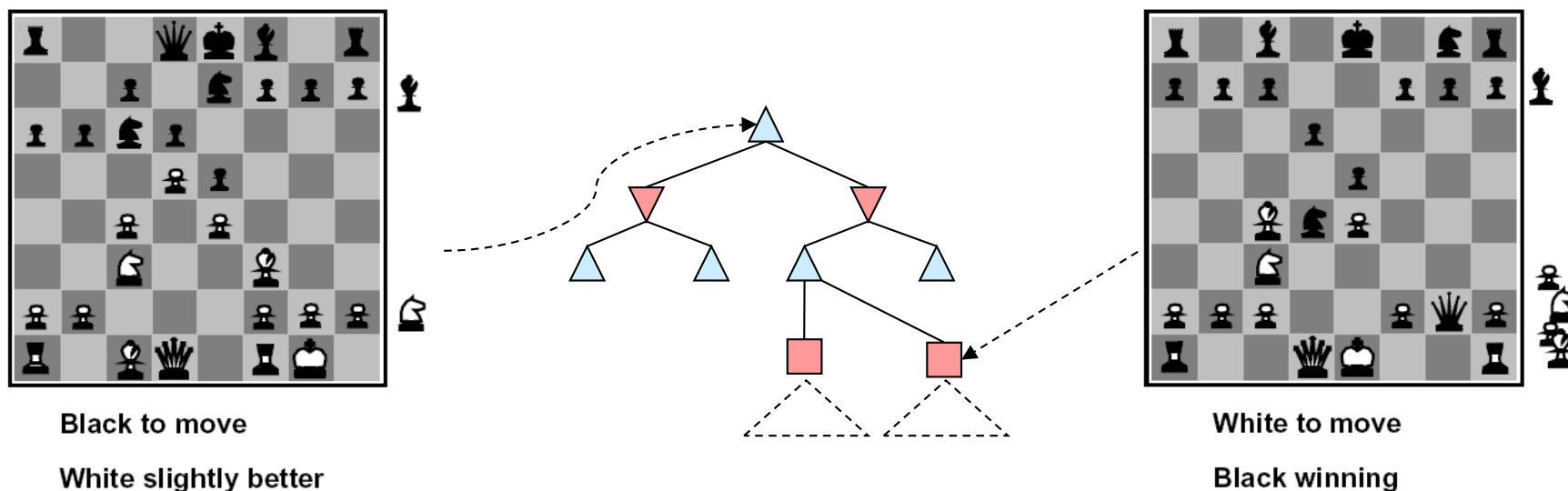


 Depth2
 Depth10





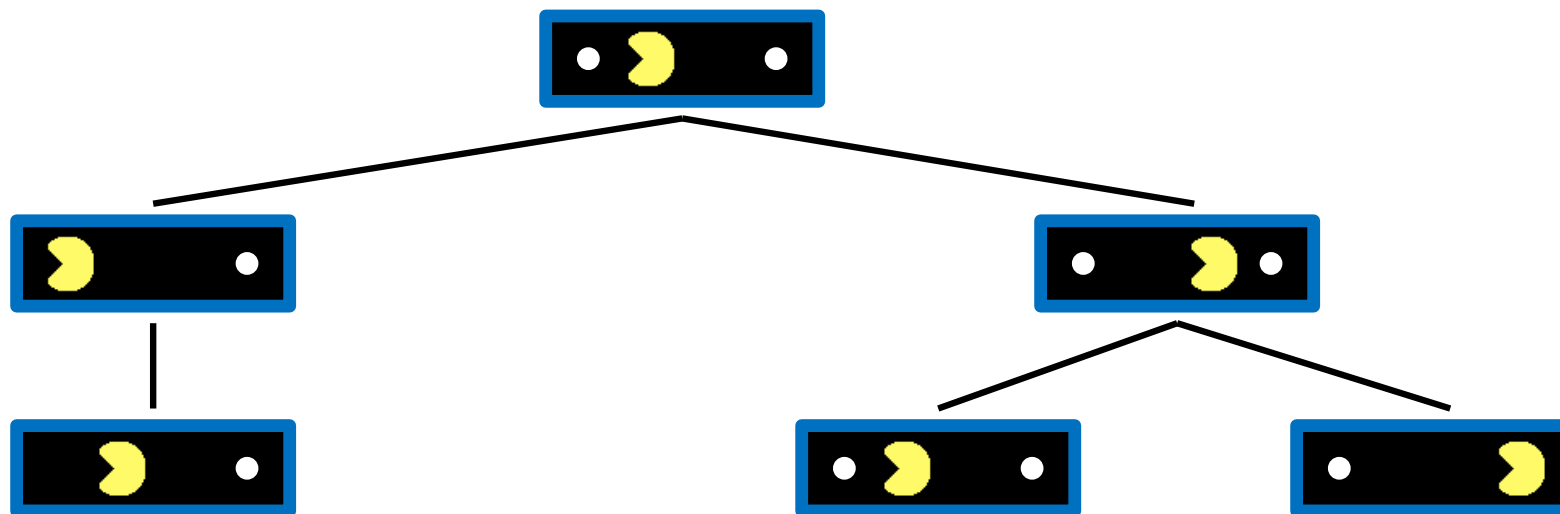
# Evaluation Functions



- Evaluation functions score non-terminals in depth-limited search
- Ideal: returns the actual minimax value of the position
- In practice: typically weighted linear sum of features:  
e.g.  $f_1(s) = (\text{num white queens} - \text{num black queens})$



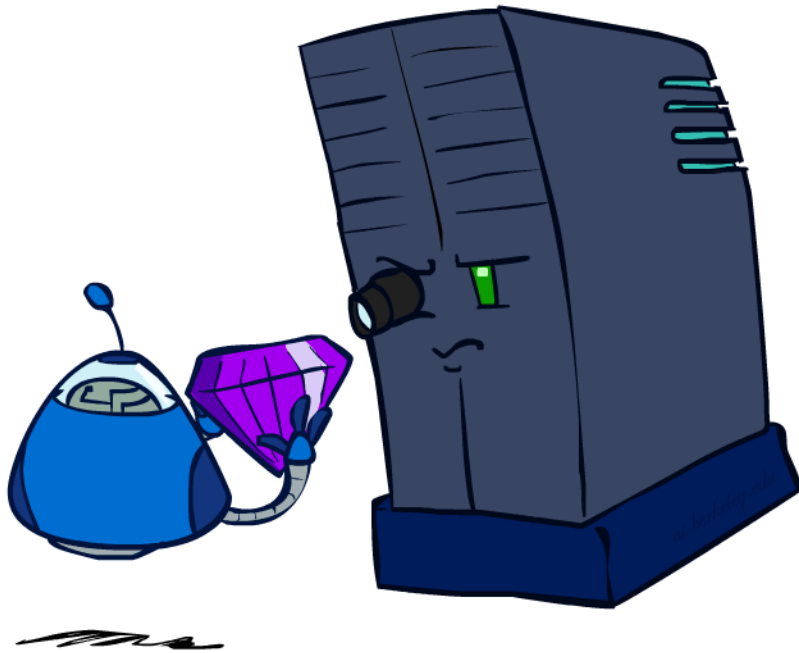
# Why Pacman Starves/Thrashes



- A danger of replanning agents!
  - He knows his score will go up by eating a dot now
  - He knows his score will go up just as much by eating a dot later
  - There are no point-scoring opportunities after eating a dot (within the horizon, two here)
  - Therefore, waiting seems just as good as eating: he may go east, then back west in the next round of replanning!



# Pacman/Ghost Evaluation



Thrashing

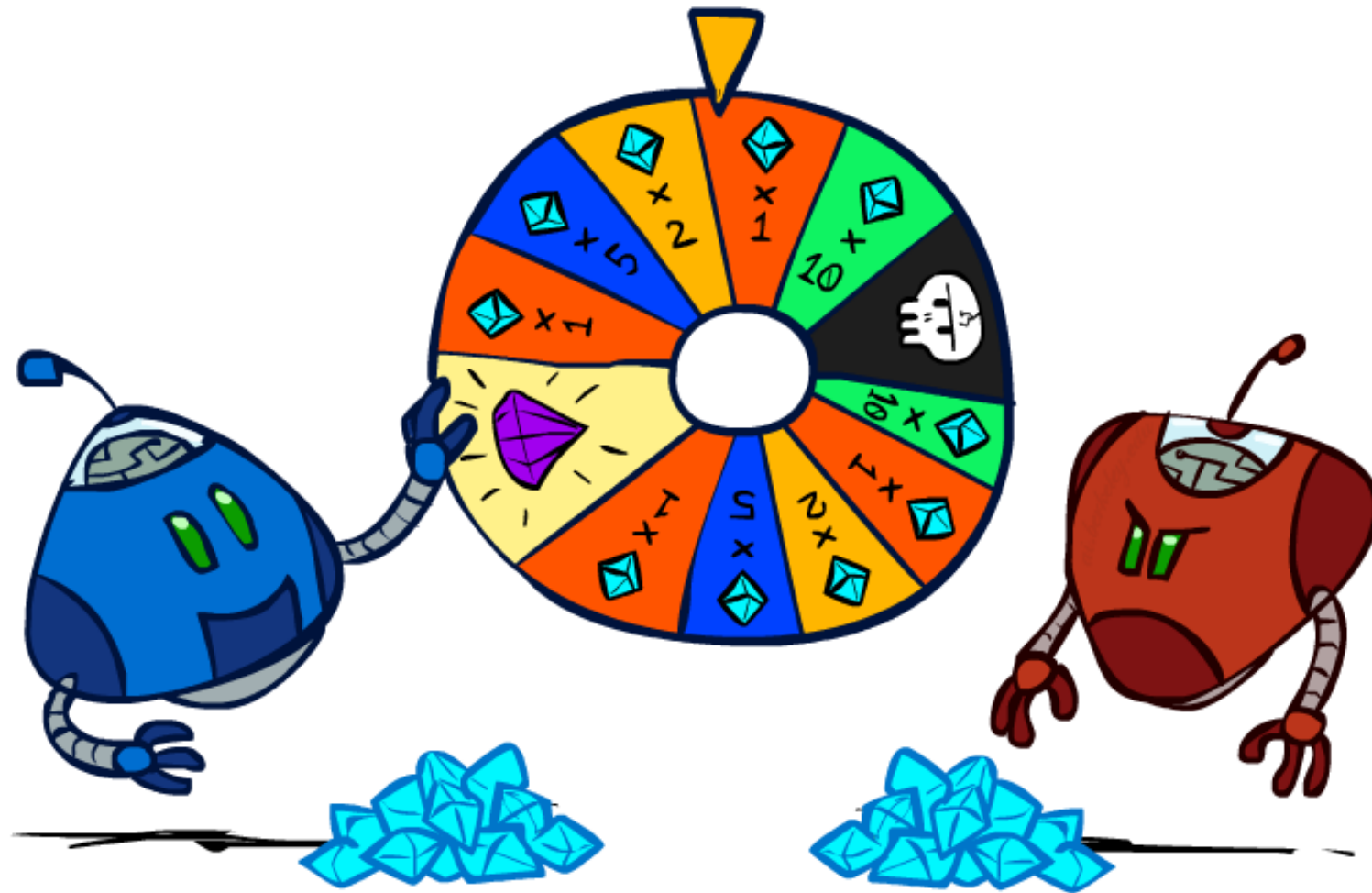
Thrashing-Fixed

SmartGhosts-1

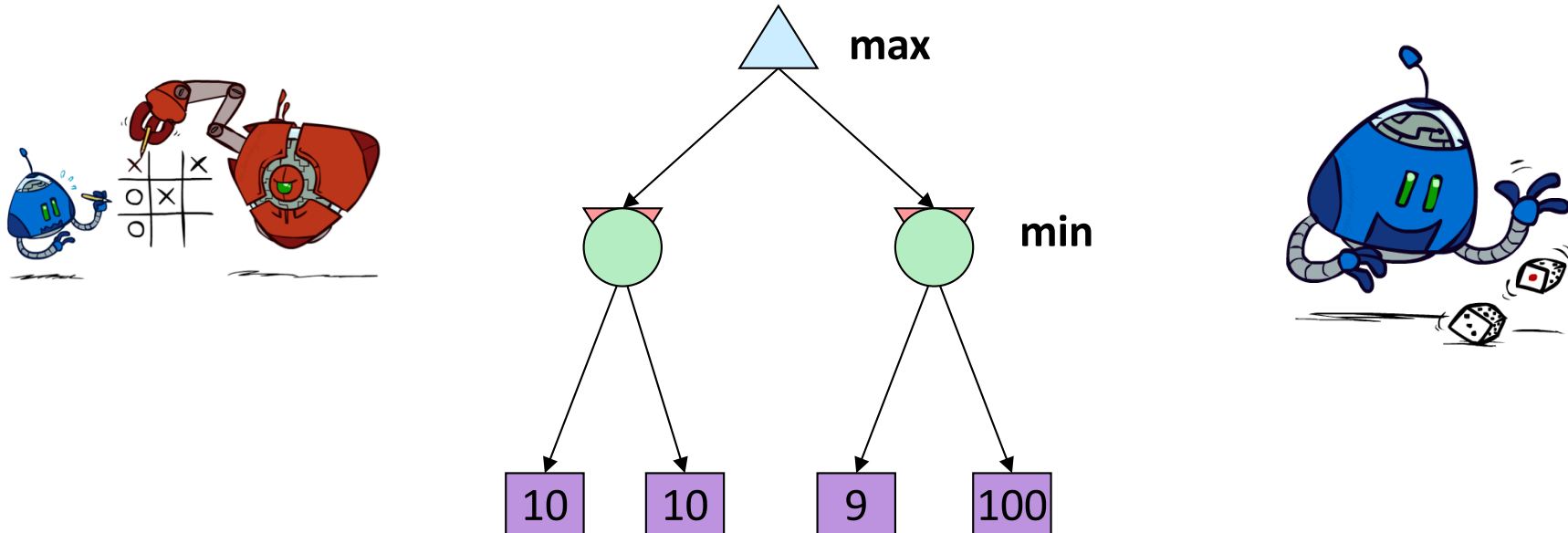
SmartGhosts-2



# Nondeterministic Games



# Worst Case vs. Average Case

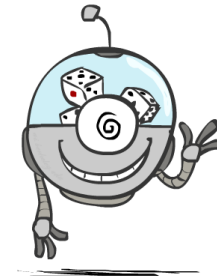
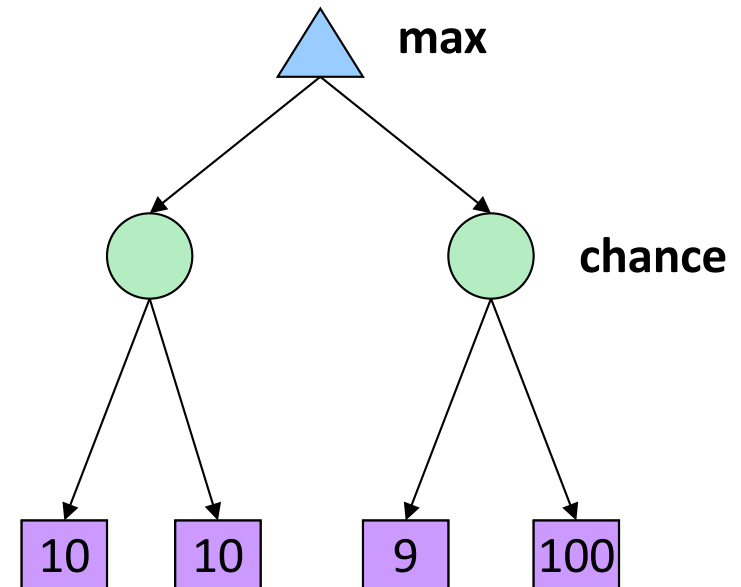


In nondeterministic games, chance is introduced by non-opponent stochasticity (e.g. dice, card-shuffling)



# Expectiminimax Search

- Why wouldn't we know what the result of an action will be?
  - Explicit randomness: rolling dice
  - Unpredictable opponents: the ghosts respond randomly
  - Actions can fail: when moving a robot, wheels might slip
- Values should now reflect average-case (**expectimax**) outcomes, not worst-case (**minimax**) outcomes
- Expectiminimax search: compute the average score under optimal play
  - Max nodes as in minimax search
  - **Chance nodes** are like min nodes but the outcome is uncertain
  - Calculate their **expected utilities**



# Reminder: Probabilities

- A **random variable** represents an event whose outcome is unknown
- A **probability distribution** is an assignment of weights to outcomes
- Example: Traffic on freeway
  - Random variable:
    - $T$  = whether there's traffic
  - Outcomes:
    - $T$  in {none, light, heavy}
  - Distribution:
    - $P(T=\text{none}) = 0.25$
    - $P(T=\text{light}) = 0.50$
    - $P(T=\text{heavy}) = 0.25$



0.25



0.50

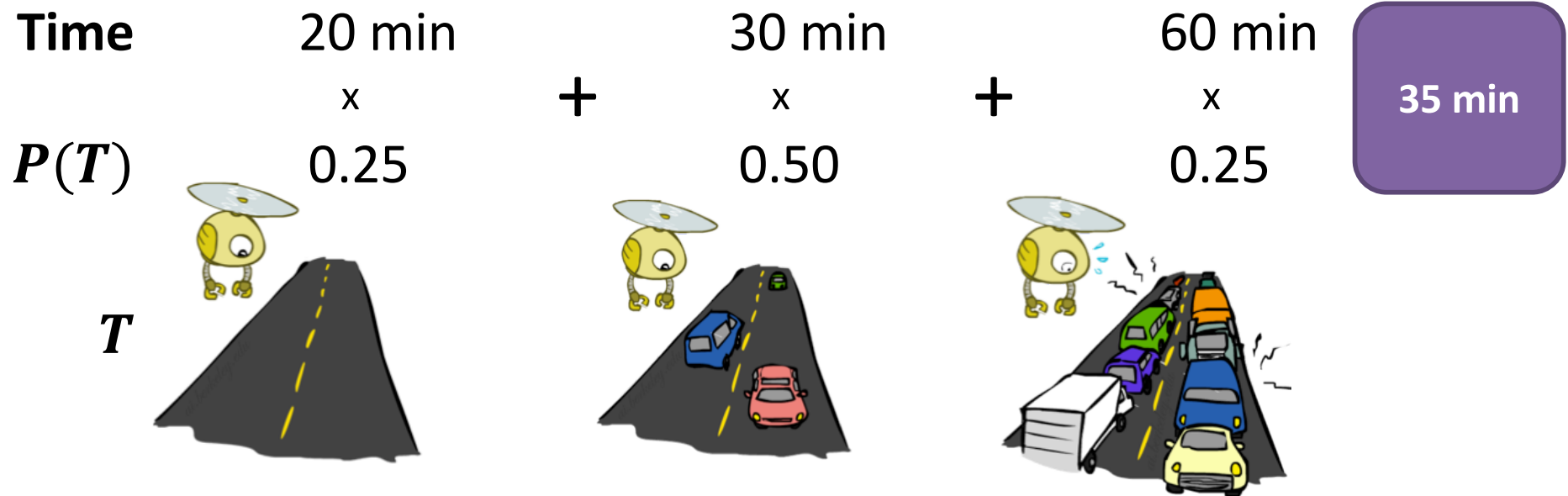


0.25



# Reminder: Expectations

- The **expected value** of a function of a random variable is the average, weighted by the probability distribution over outcomes
- Example: How long to get to the airport?





# Expectiminimax Implementation

```
def value(state):  
    if the state is a terminal state: return the state's utility  
    if the next agent is MAX: return max-value(state)  
    if the next agent is EXP: return exp-value(state)
```



```
def max-value(state):  
    initialize  $v = -\infty$   
    for each successor of state:  
         $v = \max(v, \text{value}(\text{successor}))$   
    return  $v$ 
```

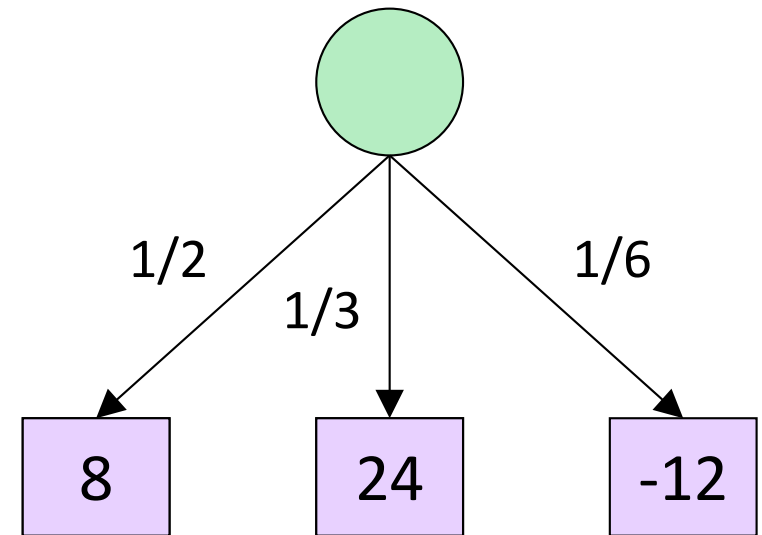


```
def exp-value(state):  
    initialize  $v = 0$   
    for each successor of state:  
         $p = \text{probability}(\text{successor})$   
         $v += p * \text{value}(\text{successor})$   
    return  $v$ 
```



# Expectiminimax Example

```
def exp-value(state):  
    initialize v = 0  
    for each successor of state:  
        p = probability(successor)  
        v += p * value(successor)  
    return v
```

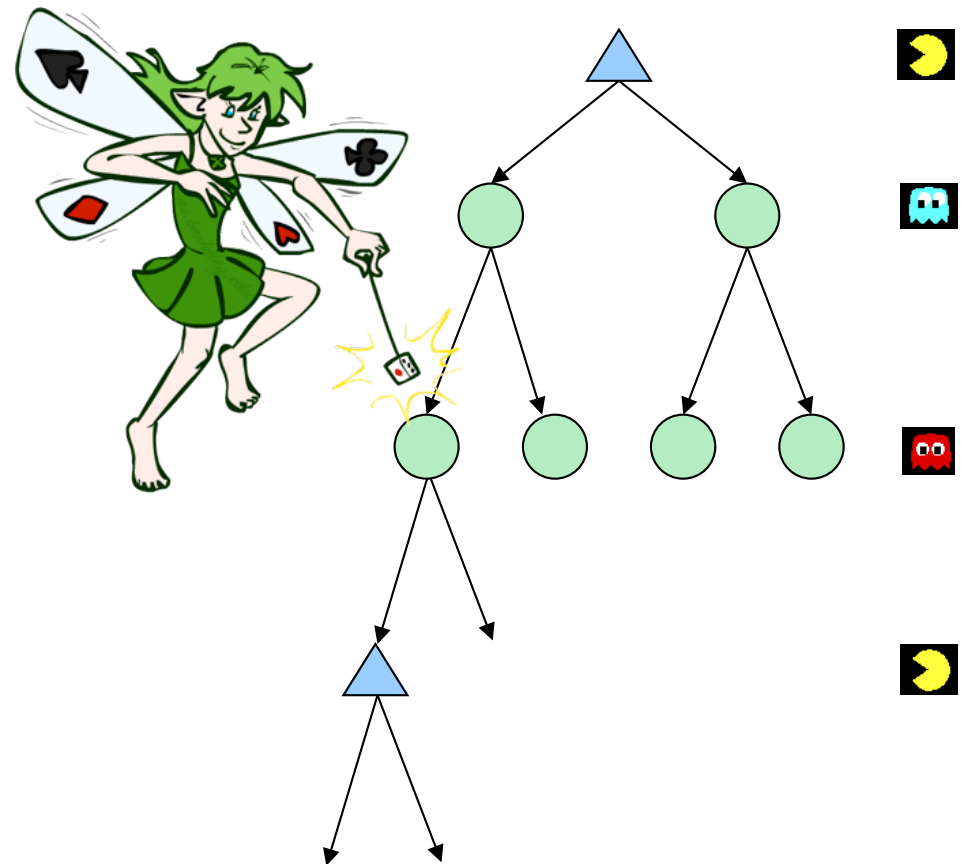


$$v = (1/2) (8) + (1/3) (24) + (1/6) (-12) = 10$$



# Where Do Probabilities Come From?

- In expectiminimax search, we have a probabilistic model of how the opponent (or environment) will behave in any state
  - Model could be a simple uniform distribution (roll a die)
  - Model could be sophisticated and require a great deal of computation
  - We have a chance node for any outcome out of our control: opponent or environment
  - The model might say that adversarial actions are likely!
- For now, assume each chance node magically comes along with probabilities that specify the distribution over its outcomes



# Summary

- A game can be formulated as a search problem, with a solution **policy** ( $S \rightarrow A$ )
- For deterministic games, the **minimax** algorithm plays optimally (assuming the **game tree** is reasonable)
- To help with resource limitations, standard practice is to employ **alpha-beta pruning** and **depth-limited search** (with an **evaluation function**)
- To model uncertainty, the **expectiminimax** algorithm introduces **chance nodes** that employ a **probability distribution** over actions to model **expected utility**

