

# Uninformed Search

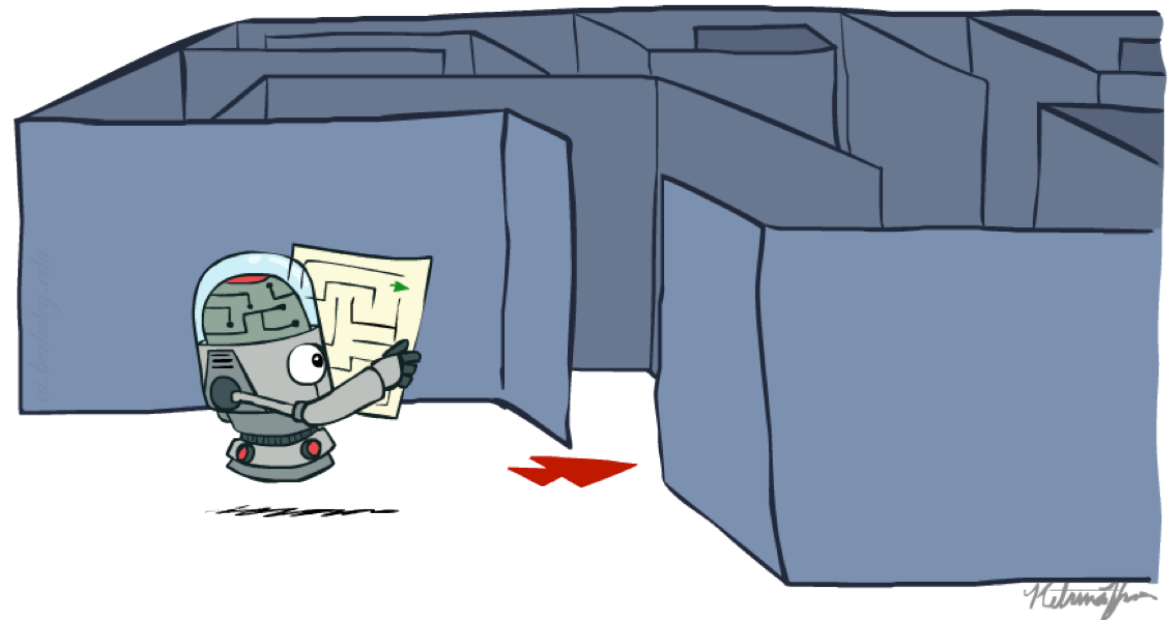
## Lecture 4

What are common search strategies that operate only given the search problem formalism? How do they compare?



# Agenda

- A quick refresher
- DFS, BFS, ID-DFS, UCS
- Unification!



# Search Problem Formalism

Defined via the following components:

- The **initial state** the agent starts in
- A **successor/transition function**
  - $S(x) = \{\text{action+cost} \rightarrow \text{state}\}$
- A **goal test**, which determines whether a given state is a goal state
- A **path cost** that assigns a numeric cost to each path

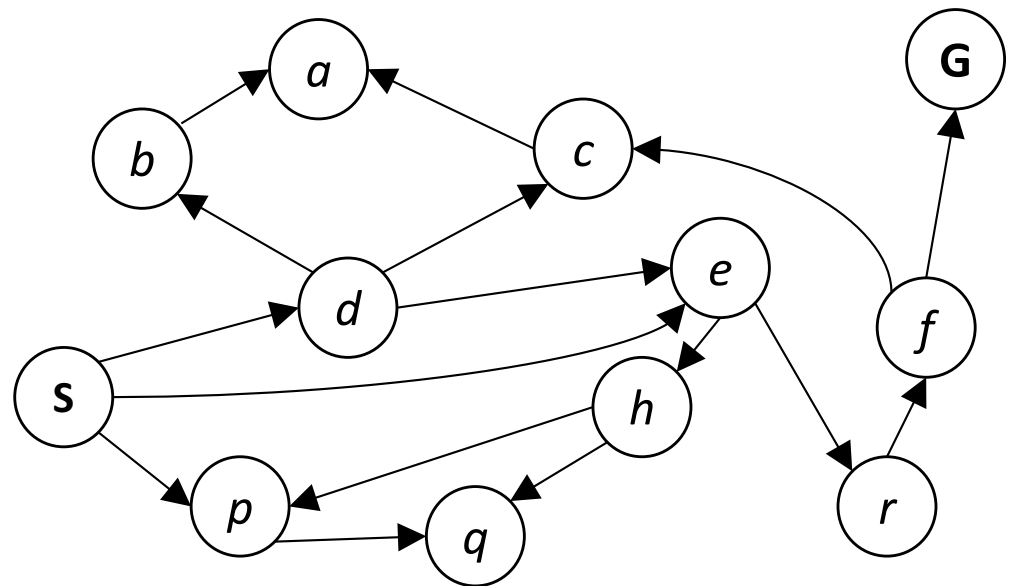
A **solution** is a sequence of actions leading from initial state to a goal state. (**Optimal** = lowest path cost.)

Together the initial state and successor function implicitly define the **state space**, the set of all reachable states



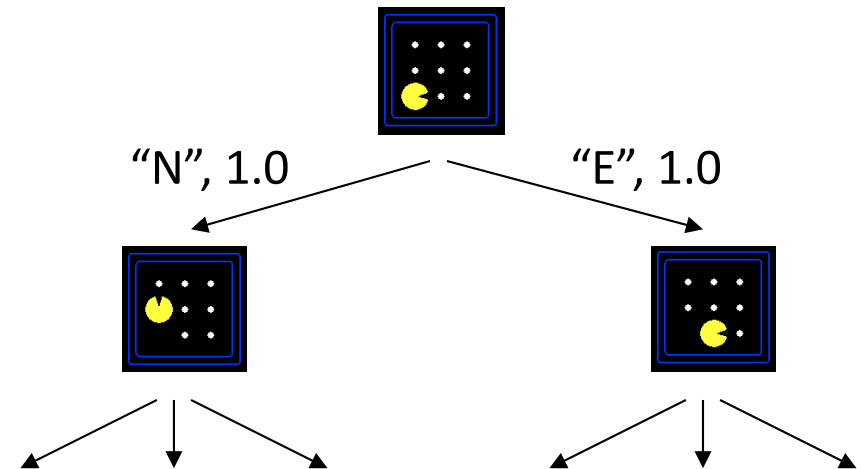
# State Space Graph

- State space graph: A mathematical representation of a search problem
  - Nodes are (abstracted) world configurations
  - Arcs represent successors (action results)
  - The goal test is a set of goal node(s)
- In a search graph, each state occurs only once!
- **We can rarely build this full graph in memory (it's too big), but it's a useful idea**



# Search Tree

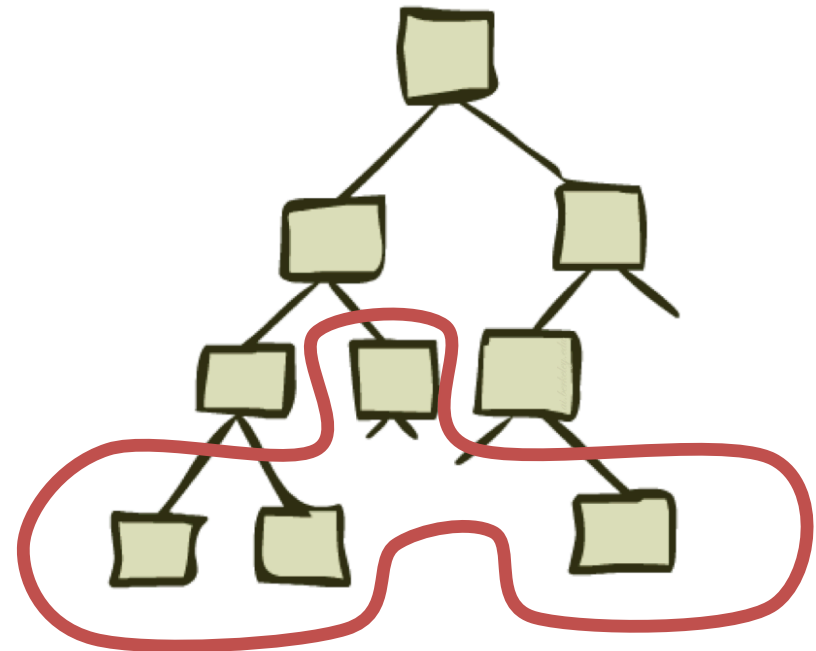
- A “what if” tree of plans and their outcomes
- The start state is the root node
- Children correspond to successors
- Nodes show states, but correspond to PLANS that achieve those states
- **For most problems, we can never actually build the whole tree**



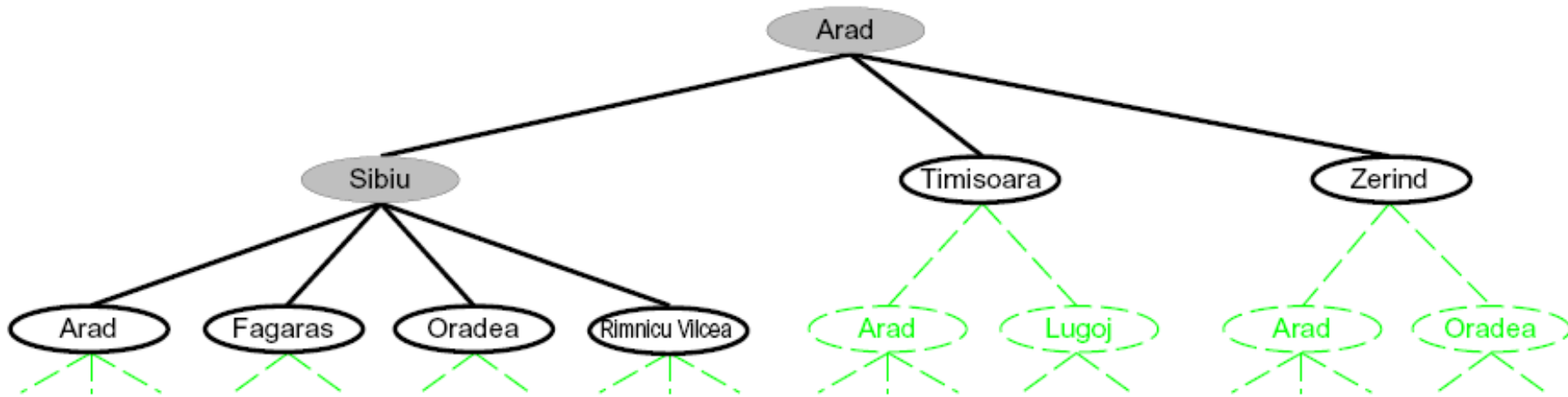
# Searching for Solutions

Basic idea: incrementally build a search tree until a goal state is found

- Root = initial state
- Expand via transition function to create new nodes
- Nodes that haven't been expanded are **leaf nodes** and form the **frontier (open list)**
- Different **search strategies** (next lecture) choose next node to expand (as few as possible!)
- Use a **closed list** to prevent expanding the same state more than once



# General Algorithm



**function** GRAPH-SEARCH(*problem*, *fringe*) **returns** a solution, or failure

*closed* ← an empty set

*fringe* ← INSERT(MAKE-NODE(INITIAL-STATE[*problem*]), *fringe*)

**loop do**

**if** *fringe* is empty **then return** failure

*node* ← REMOVE-FRONT(*fringe*)

**if** GOAL-TEST(*problem*, STATE[*node*]) **then return** *node*

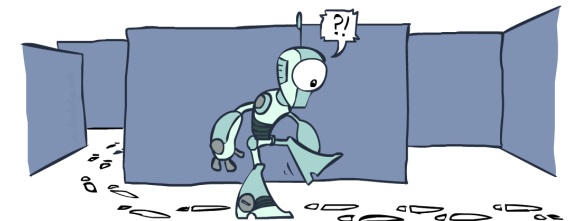
**if** STATE[*node*] is not in *closed* **then**

add STATE[*node*] to *closed*

*fringe* ← INSERTALL(EXPAND(*node*, *problem*), *fringe*)

**end**

**Queue (FIFO)**  
**Stack (LIFO)**  
**Priority Queue**



# Evaluating a Search Strategy

## Solution

- **Completeness:** does it always find *a* solution if one exists?
- **Optimality:** does it always find a least-cost solution?

## Efficiency

- **Time Complexity:**  
number of nodes generated/expanded
- **Space Complexity:**  
maximum number of nodes in memory





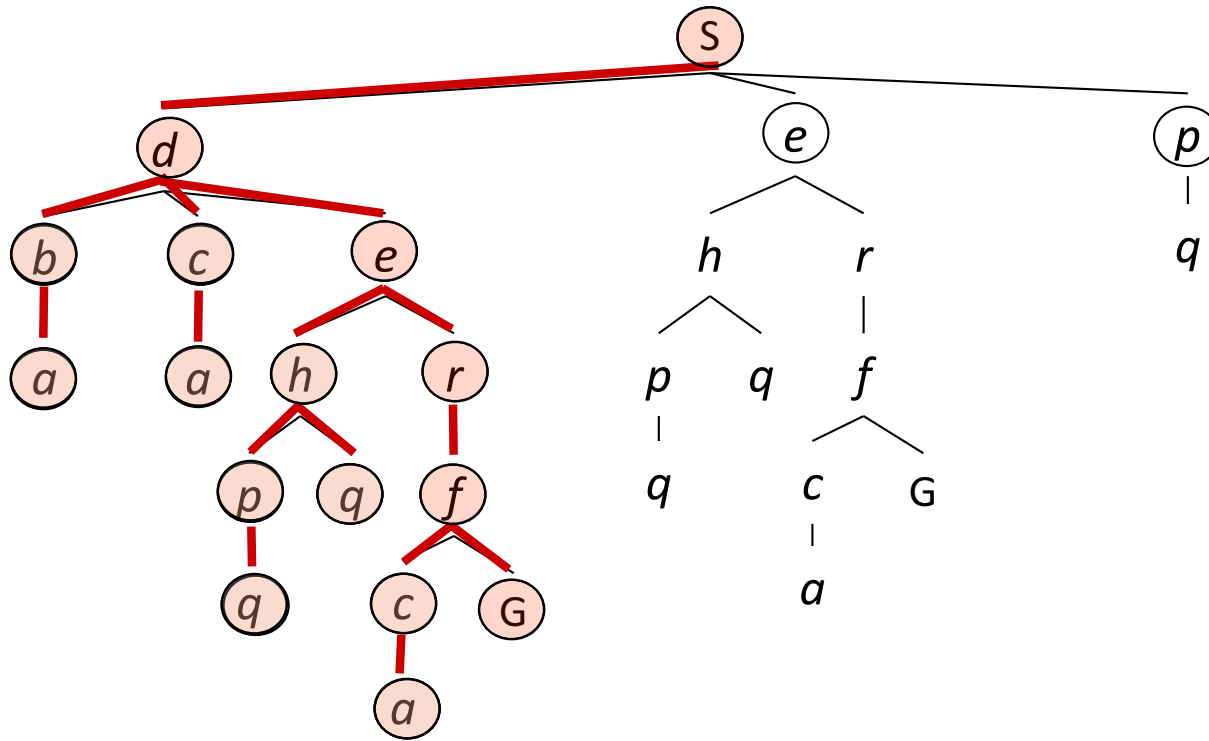
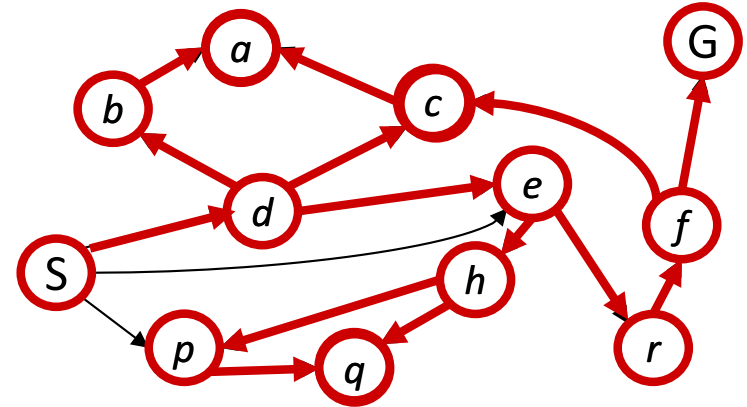
# Depth-First Search (DFS)



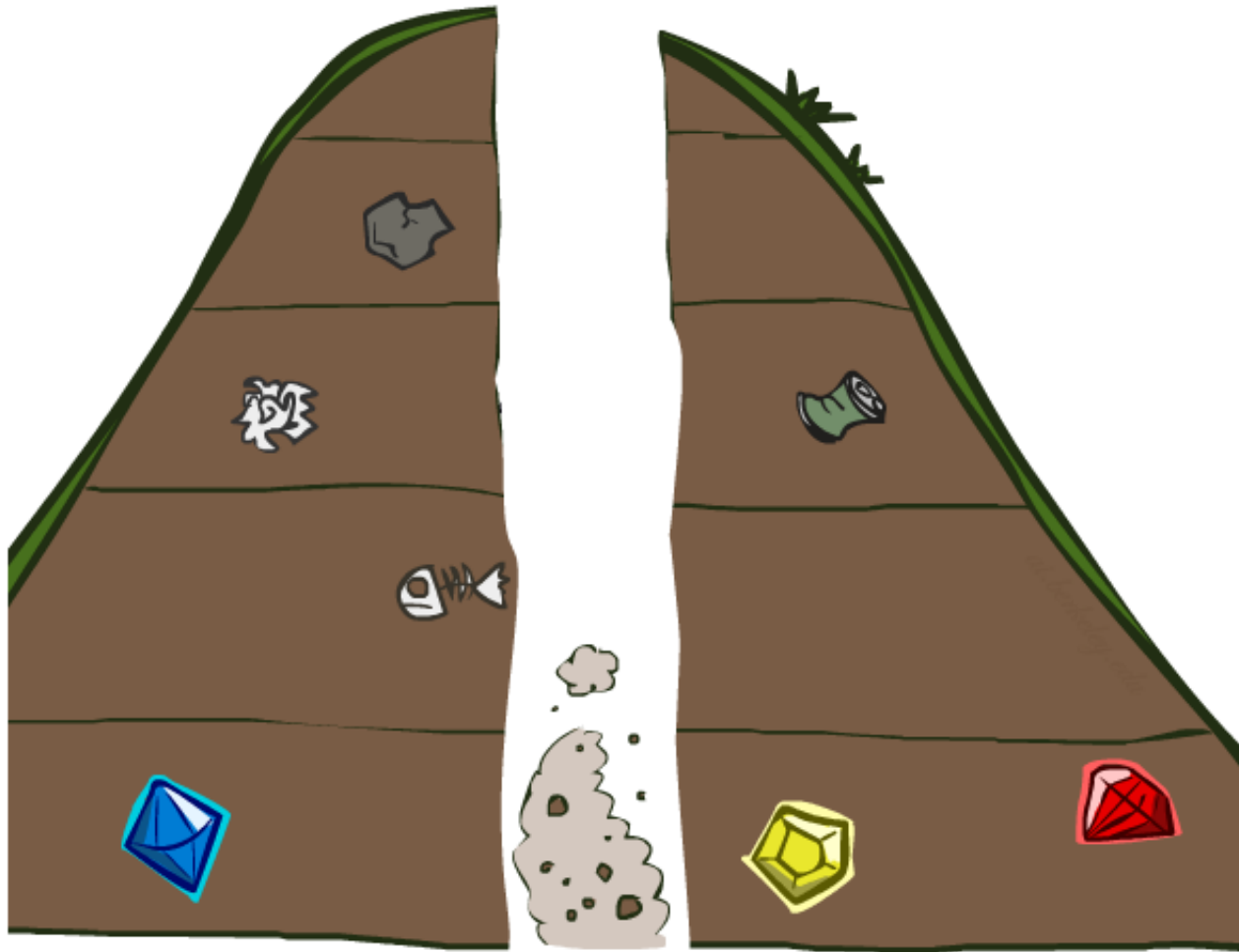
# DFS Example

Strategy: expand a deepest node first

Implementation: Fringe is a LIFO stack



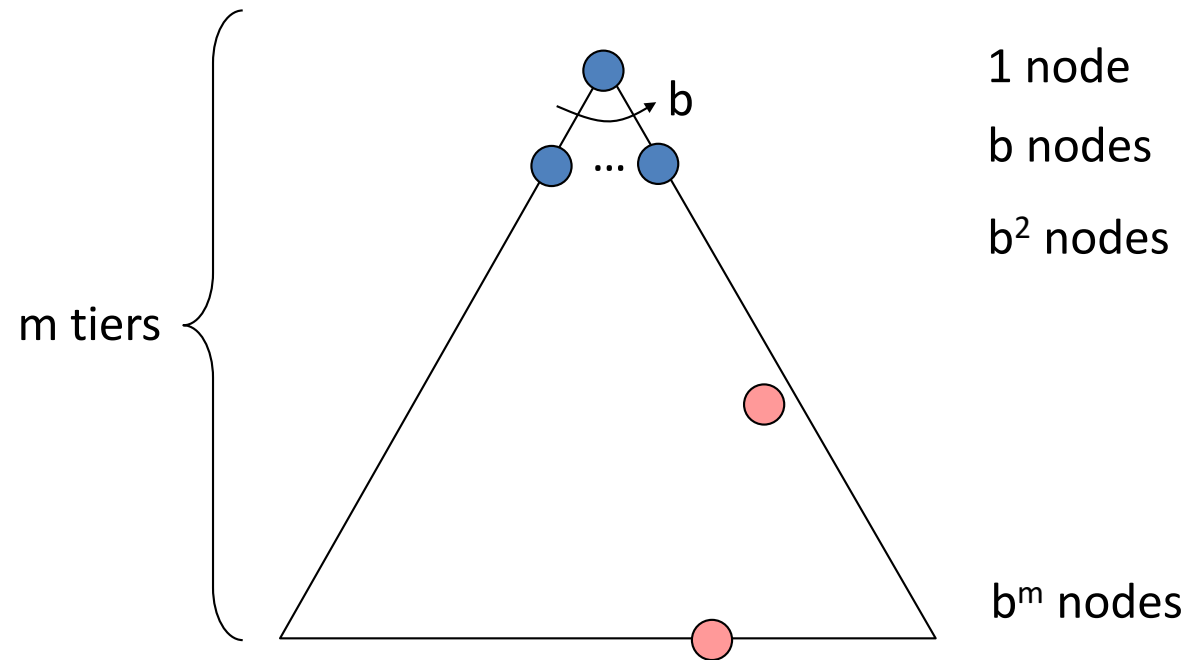
# Let's Evaluate!



# Search Tree

## Properties

- **B**ranching factor
- **M**aximum depth
- **S**olutions at various depths



Number of nodes in the tree?

- $1 + b + b^2 + \dots + b^m = O(b^m)$



# DFS Evaluation

## Time

- Expands left
  - Could be whole tree!
- Assuming finite depth,  $O(b^m)$

## Space

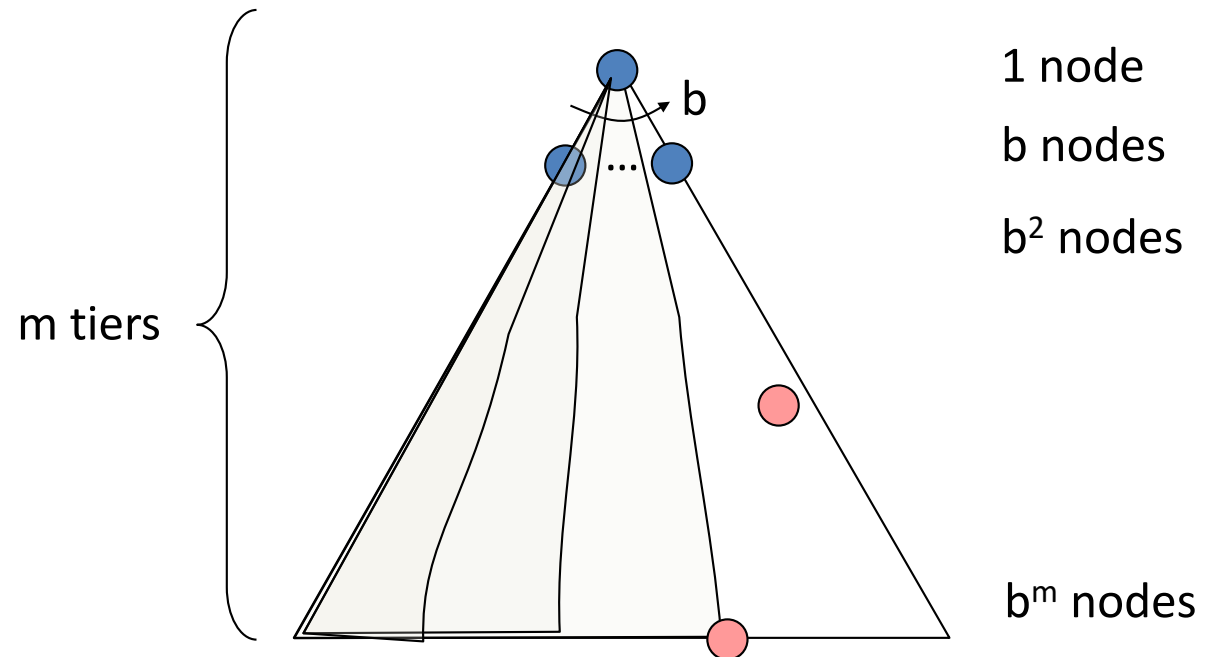
- Only siblings on path,  $O(bm)$

## Complete

- Only if finite

## Optimal

- No, "left-most" w/o regard to cost/depth



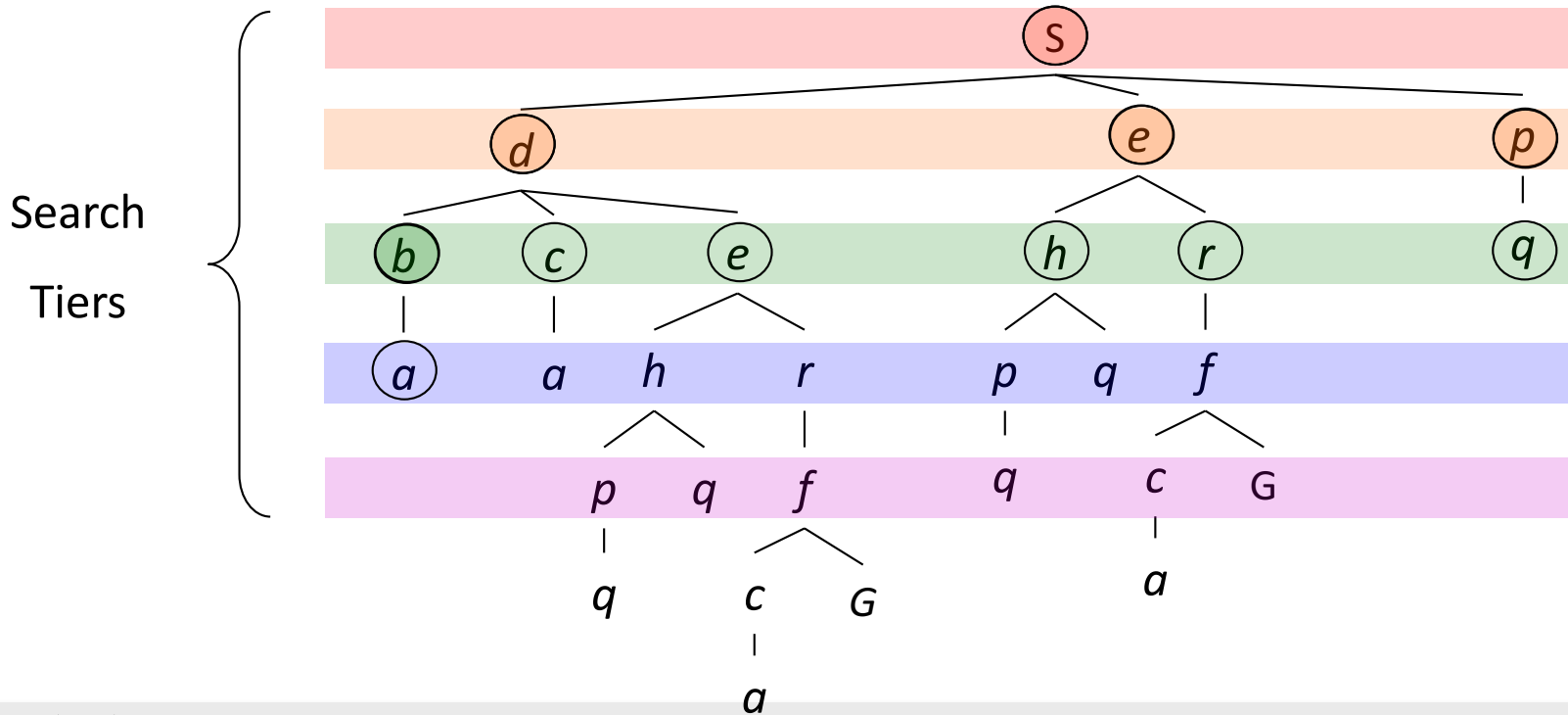
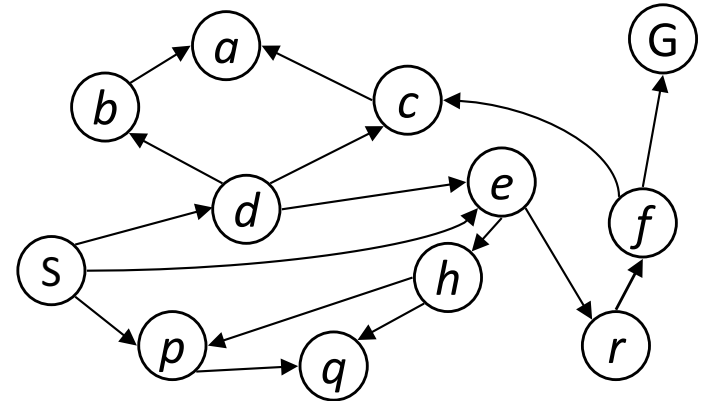
# Breadth-First Search (BFS)



# BFS Example

Strategy: expand a shallowest node first

Implementation: Fringe is a FIFO queue



# BFS Evaluation

## Time

- Processes all nodes above shallowest solution,  $O(b^s)$

## Space

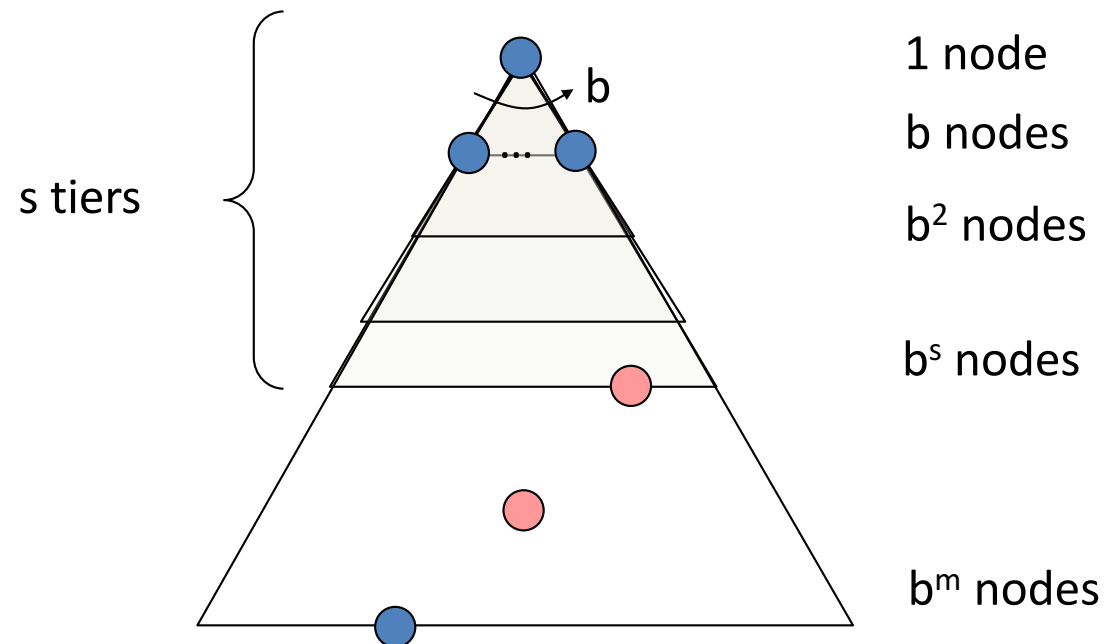
- Has roughly the last tier, so  $O(b^s)$

## Complete

- Yes!

## Optimal

- Only if all costs equal (more later)





# DFS vs. BFS

 Empty-DFS

 Empty-BFS

 Maze-DFS

 Maze-BFS



# Grounding the Branching Factor

Depth	Nodes	Time	Memory
2	110	0.11 msec	107 KB
4	11,110	11 msec	10.6 MB
6	$10^6$	1.1 sec	1 GB
8	$10^8$	2 mins	103 GB
10	$10^{10}$	3 hours	10 TB
12	$10^{12}$	13 days	1 PB
14	$10^{14}$	3.5 years	99 PB
16	$10^{16}$	350 years	10 EB

## Assumptions

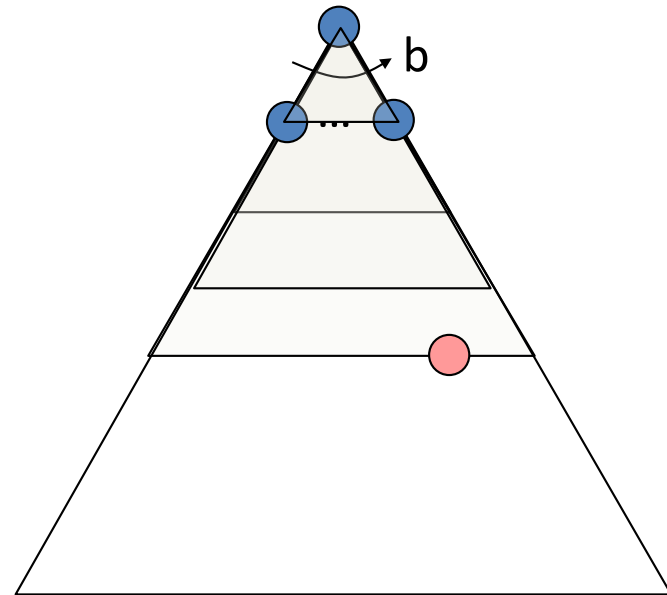
- $b = 10$
- 1 million nodes/second
- 1000 bytes/node

Memory often becomes the limiting factor

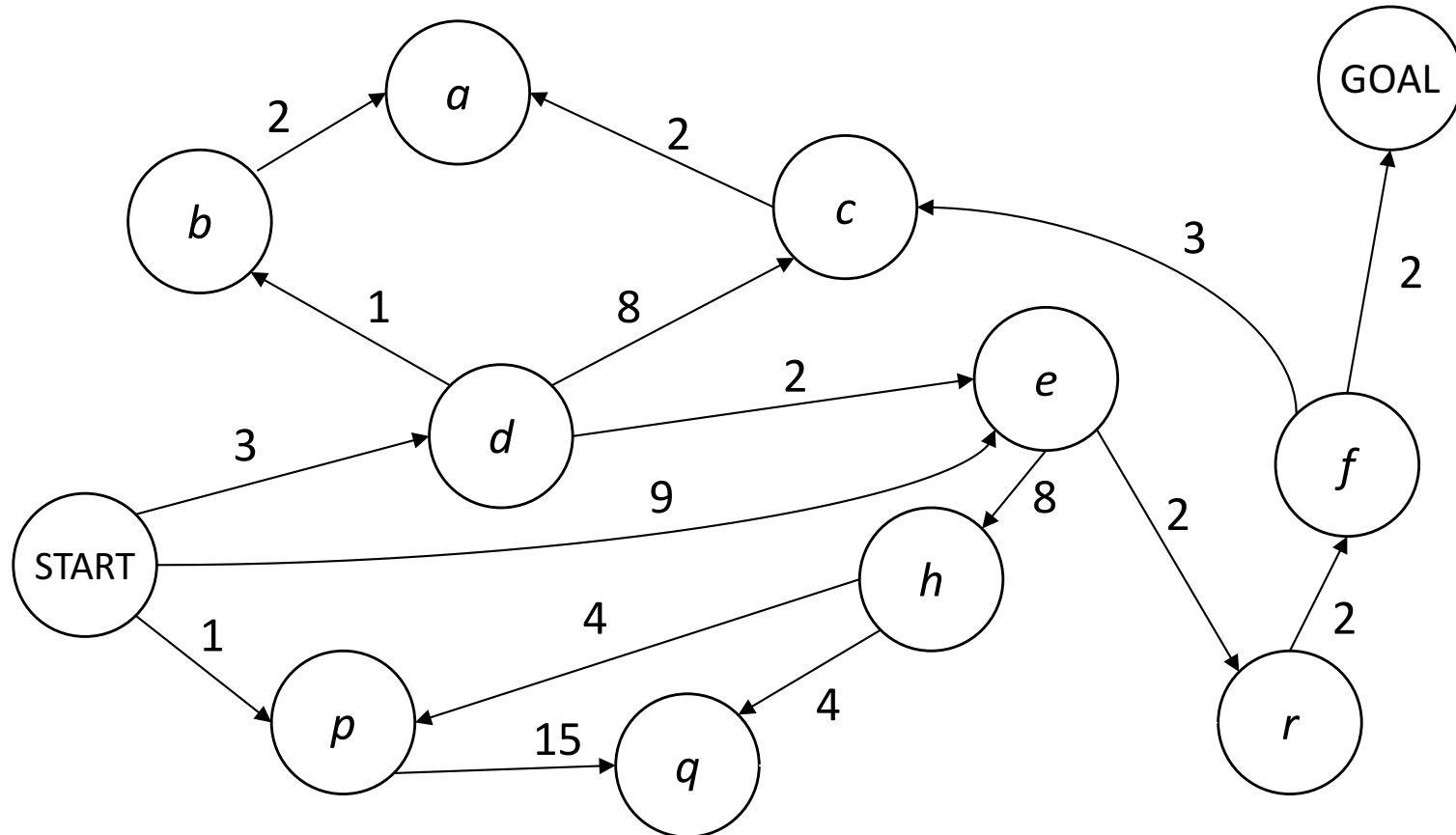


# Iterative Deepening DFS

- Basic idea: DFS memory with BFS time/shallow solution
  - DFS up to 1
  - DFS up to 2
  - ....
- Generally most work happens in the lowest level searched, so not too wasteful



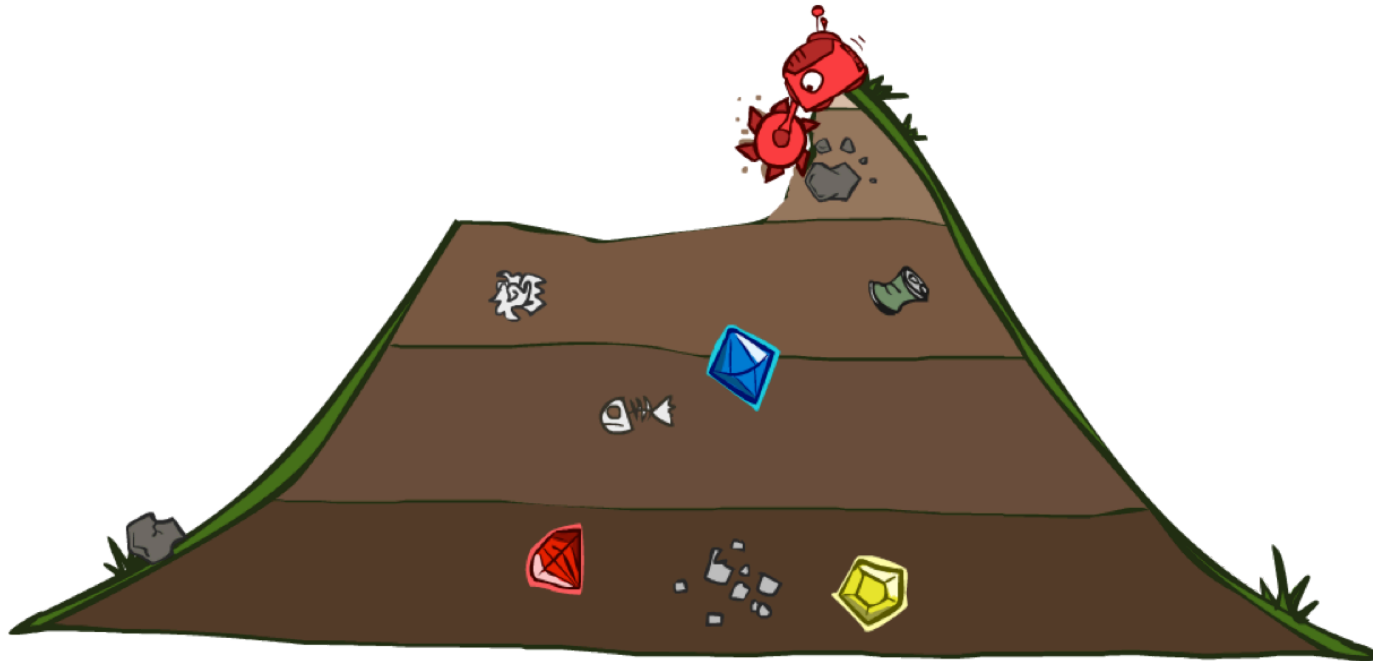
# Cost-Sensitive Search



- BFS finds the shortest path in terms of number of actions, but it does not find the least-cost path.
- We will now cover a similar algorithm which does!



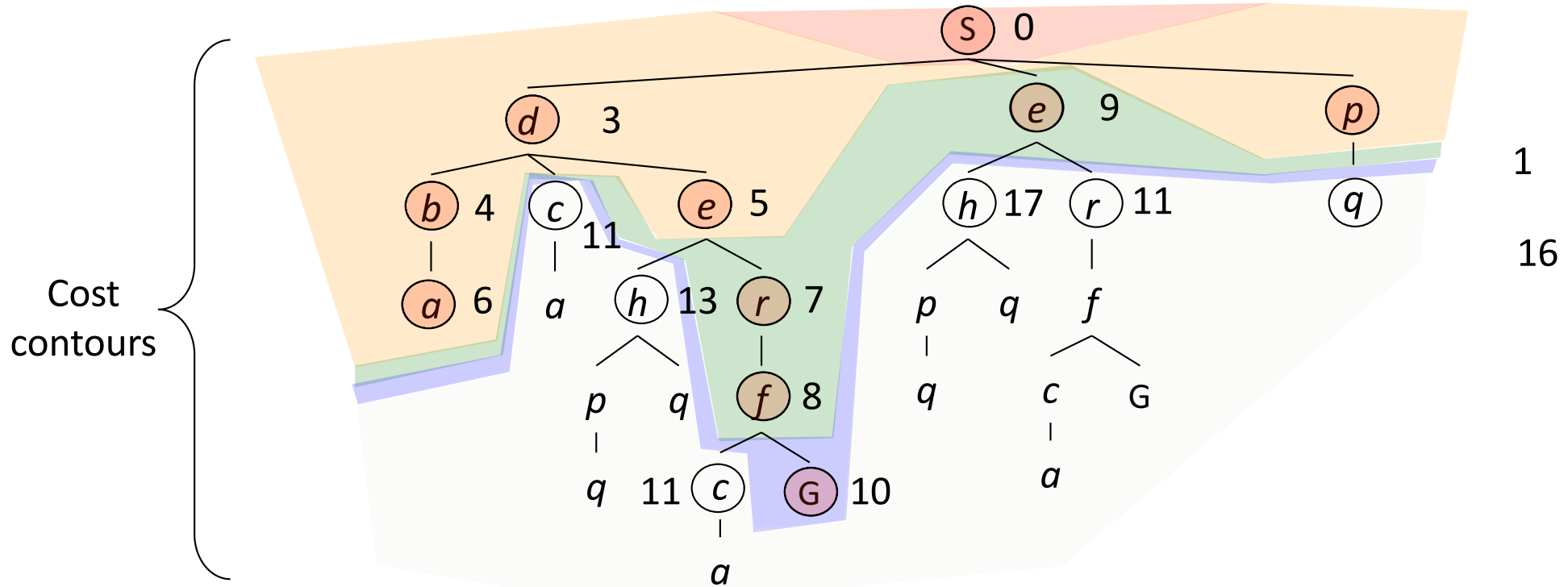
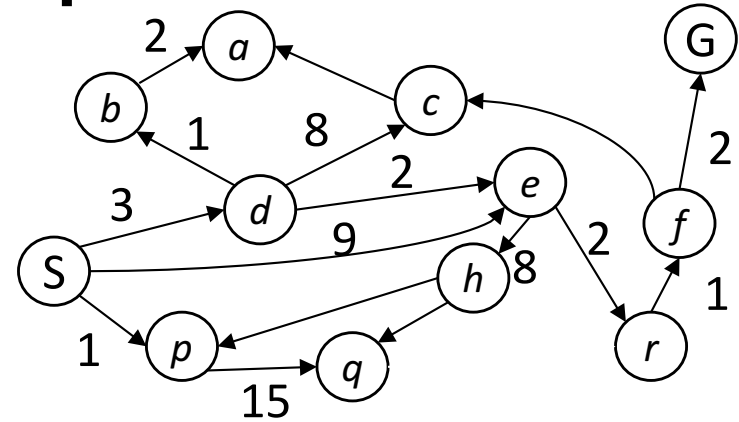
# Uniform-Cost Search (UCS)



# UCS Example

Strategy: expand a cheapest node first

Fringe is a priority queue  
(priority: cumulative cost)



Cost contours



# UCS Evaluation

## Time

- $O(b^{C^*/\epsilon})$

## Space

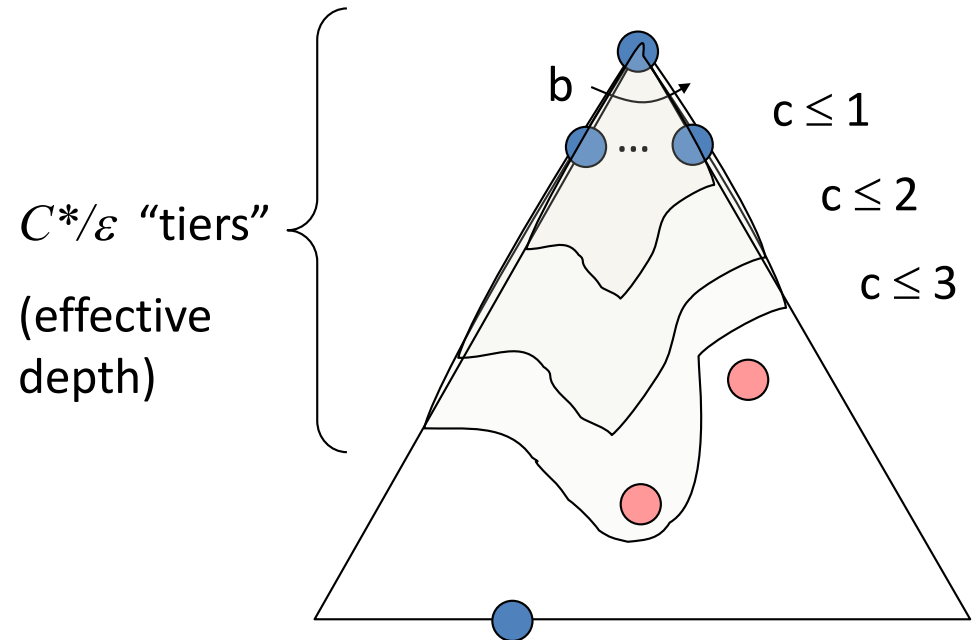
- $O(b^{C^*/\epsilon})$

## Complete

- Yes!

## Optimal

- Yes!

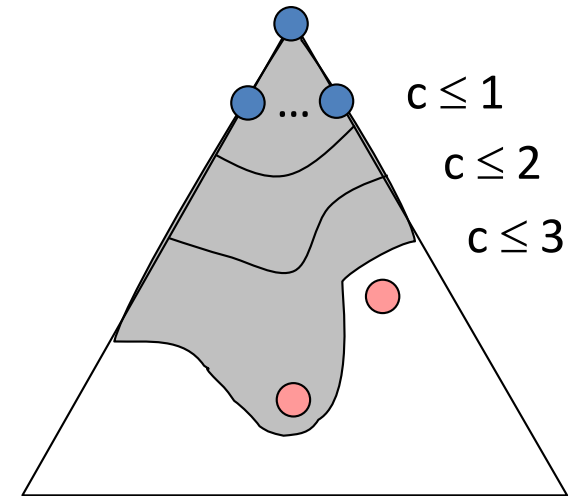
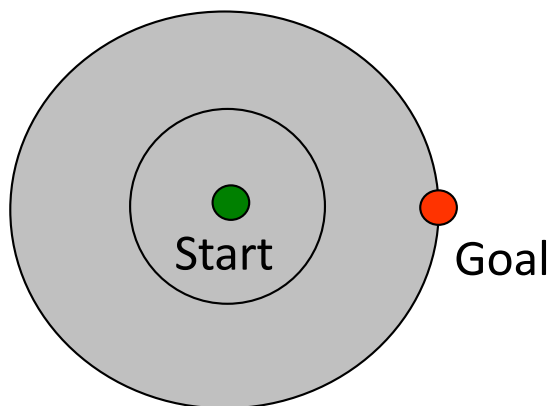







Assume solution costs  $C^*$  and arcs cost at least  $\epsilon$



# UCS vs. DFS vs. BFS

- UCS is good and optimal
- However, it still moves in every direction – it's not **informed** about goal direction...



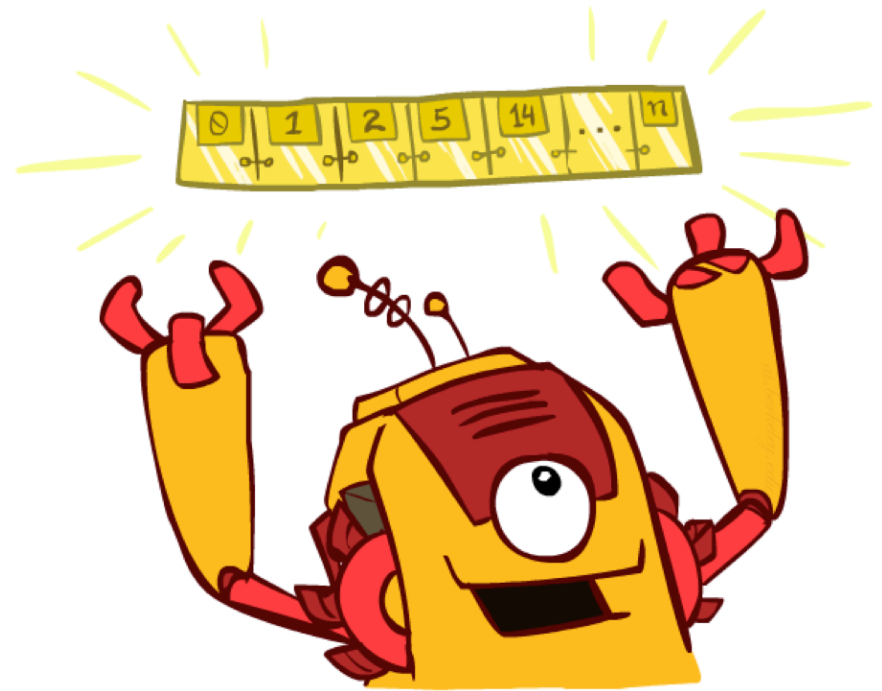
	Empty-UCS
	Maze-UCS
	MazeCost-DFS
	MazeCost-BFS
	MazeCost-UCS





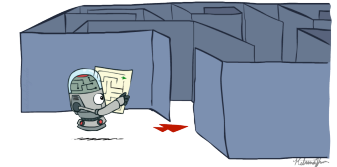
# Unification

- All these search algorithms are the **same** except for fringe strategies
- Conceptually, all fringes are priority queues (i.e. collections of nodes with attached priorities)
- Practically, for DFS and BFS, you can avoid the  $\log(n)$  overhead from an actual priority queue, by using stacks and queues

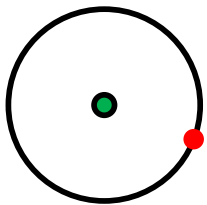


# Uninformed Search

Search given **only** the problem definition

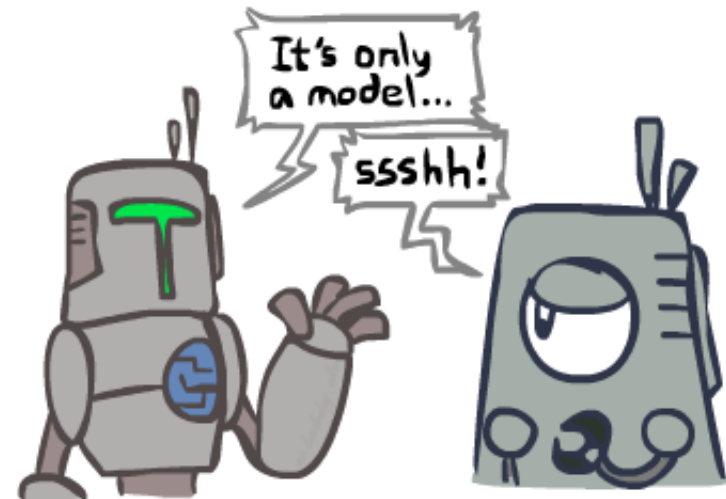


	DFS	BFS	UCS
Fringe	LIFO (stack)	FIFO (queue)	PQ (path cost)
Complete		X	X
Optimal			X
Time	$O(b^m)$	$O(b^s)$	$O(b^{C^*/\epsilon})$
Space	$O(bm)$	$O(b^s)$	$O(b^{C^*/\epsilon})$
Assumptions: potentially infinite depth, arbitrary positive action costs			



# A Reminder

- Search operates over models of the world
- The agent doesn't actually try all the plans out in the real world!
- Planning is all “in simulation”
- Your search is only as good as your models...



# Search Gone Wrong

Google Maps

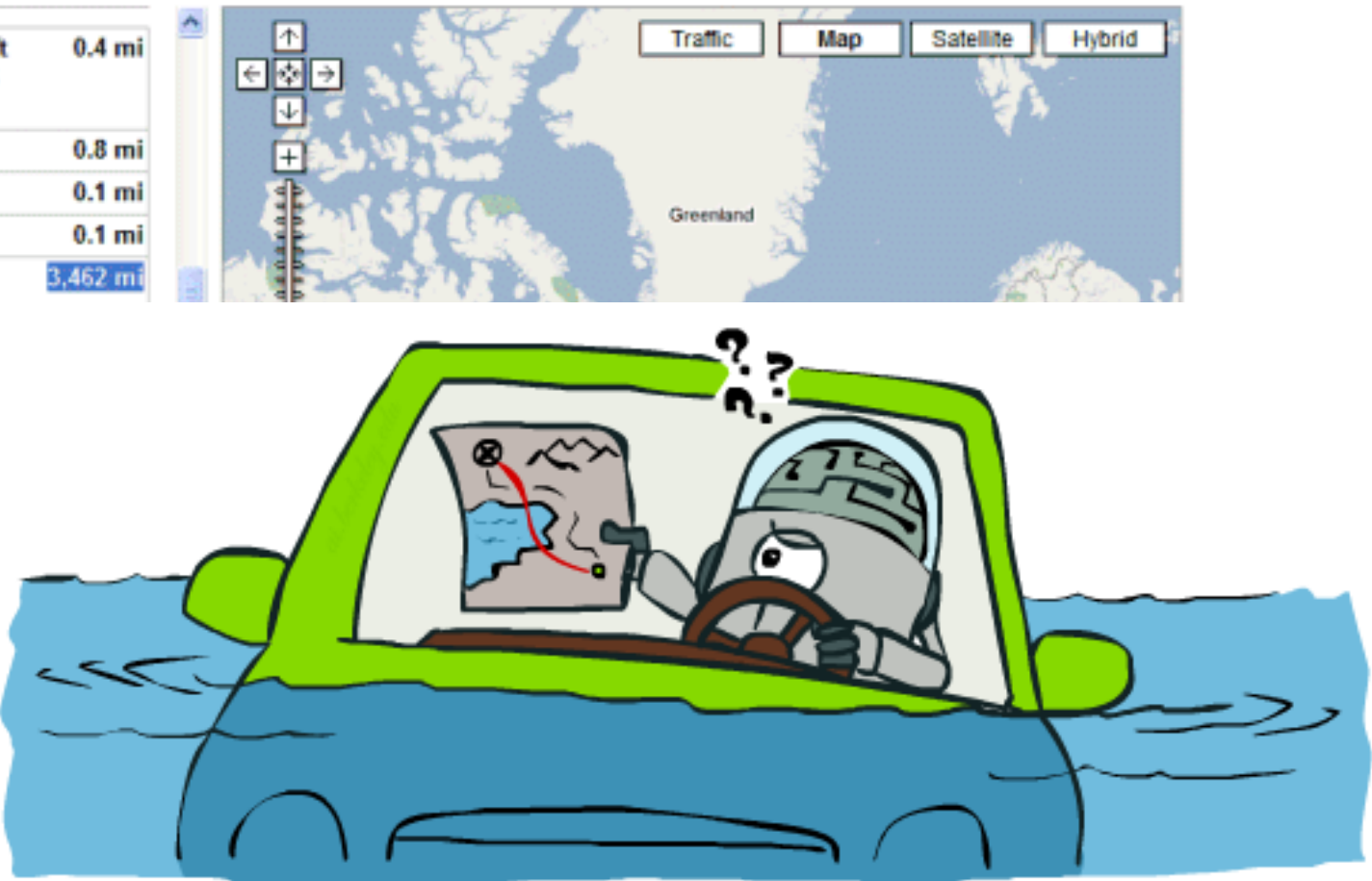
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29.	Take exit 24 A-B-C on the left toward I-93 N/Concord NH/S Station/I-93 S/Quincy	0.4 mi
30.	Merge onto Atlantic Ave	0.8 mi
31.	Turn right at Central St	0.1 mi
32.	Turn right at Long Wharf	0.1 mi
33.	Swim across the Atlantic Ocean	3,462 mi
34.	Slight right at E05	
35.	At the roundabout, take the exit onto E05/Pont Vauban	
36.	Turn right at E05	
37.	Take the exit onto A29/E44 toward Amiens Toll road	
38.	Take the exit toward Dieppe Amiens/Calais/A151/Rouen	



# Summary

- We evaluated several uninformed strategies to solve a search problem
  - **DFS, ID-DFS, BFS, UCS**
- DFS, BFS, and UCS can all be implemented via a generic graph-search algorithm over a search tree by simply changing how the fringe is organized

