#### Indexes

#### Lecture 9



#### Outline

- Context
- Functionality
- Index types
- Utility
- Tradeoffs and considerations
  - Selectivity
- A reminder of syntax



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#### **Database Design and Implementation Process**

<b>Figure 10.1</b> Phases of database design and implementation for large databases.		Data content, structure, and constraints			Database applications	
Phase 1:	Requirements collection and analysis	-	Data requirements		Processing requirements	
Phase 2:	Conceptual database design		Conceptual Schema design (DBMS-independent)		Transaction and application design (DBMS-independent)	
Phase 3:	Choice of DBMS		•			
Phase 4:	Data model mapping (logical design)		Logical Schema and view design (DBMS-dependent)	Frequence performa constrain	ance	
Phase 5:	Physical design		Internal Schema design (DBMS-dependent)			
Phase 6:	System implementation and tuning		DDL statements SDL statements		ransaction and application implementation	



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### What is an Index?

- Persistent data structure, stored in the database
- Primary mechanism to get improved query performance
- Many interesting issues (see Ch. 17-18); we will focus on usage, tradeoffs



#### Functionality

An index answers certain kinds of questions very efficiently (depends upon type)

- Equality: fieldname=value
- Range/ordering: fieldname>value
  - Only index that maintains ordering (e.g. tree-based)

# Can be used for WHERE clause, as well as JOIN and ORDER BY



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# Comparison (1)

SELECT * FROM T WHERE	т	Α	В	С
	1	cat	1	
<ul> <li>No indexes (indices) Anything = full table scan (slow)</li> </ul>	2	dog	3	
<ul> <li>Index on (A)</li> </ul>	3	panda	7	
A = 'panda' (fast)	4	cat	4	
<ul><li>A &gt; 'dog' (fast, if ordered)</li><li>ORDER BY A (fast, if ordered)</li></ul>	5	cat	5	
<ul> <li>Index on (B)</li> </ul>	6	panda	9	
B = 1 (fast) B >= 5 (fast, if ordered)	7	moose	10	
ORDER BY B (fast, if ordered)	8	dog	8	
<ul> <li>Index on (A, B)</li> <li>B = 1 AND A = (cat) (foot)</li> </ul>	9	dog	10	

- B = 1 AND A = 'cat' (fast)
  A >= 'cat' AND B = 3 (fast, if ordered)
  A = 'panda' ORDER BY B (fast, if ordered)
- Index on (C,A), (C,B), ... (i.e. start with C) Anything not including C = full table scan (slow)



#### Comparison (2)

T1	Α	В	С	Т2	X	Y	Z
1	cat	1		i	felidae	1	
2	dog	3		ii	canidae	3	
3	panda	7		iii	bear	7	
4	cat	4		iv	felidae	4	

#### T1 JOIN T2 ON T1.B=T2.Y

- No indexes: slow
- Index on T1(B): scan T2, fast search in T1
- Index on T2(Y): scan T1, fast search in T2
- Index on T1(B), T2(Y): merge sort (if ordered)



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# Index Types

- Clustered vs. Non-clustered
- Covering (w.r.t. a query)
- Balanced Trees (B+-Trees)
- Hash Tables



#### Clustered vs. Non-clustered

- Clustered: affects physical order on disk
  - At most one per table (for some RDBMS, PK)
  - Fast when data accessed in order/reverse
- Non-clustered: induces logical ordering
  - Arbitrary number per table
  - Typically non-prime
  - For some RDBMS, UNIQUE



### Covering

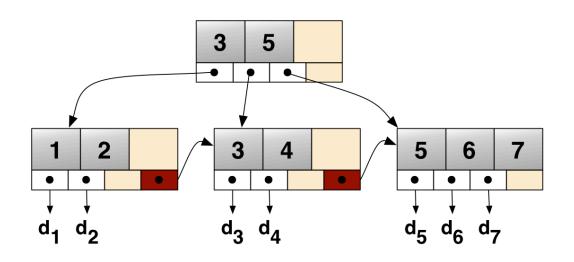
- Typically indexes help the DBMS *find* the row of interest
  - ID -> Name
  - Name->ID

ID	Name
1	Alice
2	Bob
3	Carol
4	Dan

- A covering index contains all the necessary data within the index itself (w.r.t. to query or queries)
  - More storage vs. IO savings
  - (ID, Name) or (Name, ID)



#### **B+-Trees**



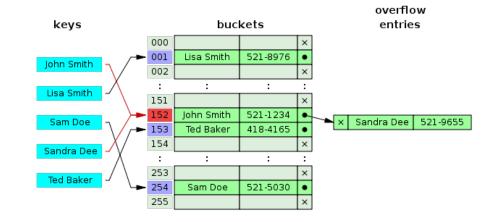
- Balanced, constant out-degree (within range)
- Values (i.e. row pointer) only at leaves
  - Distinguishes from a B-tree
  - Linked list at leaves, in order
- Logarithmic traversal, constant at leaf
- Typical default index for DBMS; also used in file systems, etc.



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#### Hash Table



- "Constant" access time (under certain assumptions, amortized)
- No range queries



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#### Utility

#### Pro

 Can make the difference between full table scan and log/constant lookup

#### Con

- Extra space
  - Linear with # rows
- Extra time
  - Creation (moderate)
  - Maintenance (can offset savings)



## Choosing the Index(es) to Create

- Table size
  - Many rows = larger cost to table scan
- Data distribution (selectivity)
  - Fewer distinct values = higher likelihood needing to touch many rows, independent of index usage
    - Index can lead to lots of IO/cache misses vs. sequential scan via clustered index
- Query vs. update load
  - Many updates = higher relative index maintenance cost



### Selectivity

- Cardinality: # distinct values in a column SELECT COUNT(DISTINCT col\_name)
   FROM table\_name;
- **Selectivity**: 100% \* cardinality / # rows
  - Compare for 10K rows...
    - Gender (M/F)
    - Country (195 + Taiwan)
    - Birthday (Jan. 1 -> Dec. 31)



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#### **General Advice**

- Use narrow indexes (i.e. few columns); these are more efficient than compound indices
- Avoid a large number of indices on a table
- Avoid "overlapping" indices that contain shared columns (often a single index can service multiple queries)
- For indices that contain more than one column: given no other constraints, place the most selective column first
- Always define one clustered index on each table (typically equates to defining a PK)



## Creating Indexes CREATE [UNIQUE] INDEX index\_name ON table\_name (c\_name1, ...) [OPTIONS];

<u>Notes</u>

- Ordering of columns is VERY important
- Options often refer to the type of index being used (e.g. btree, hash) and other important flags

